

A Hypervisor IPS based on Hardware Assisted Virtualization Technology

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Presentation Outline

- 1. Review of subversive techniques in kernel space
- 2. Review of Virtualization Technology
- 3. Viton, Hypervisor IPS
- 4. Conclusion





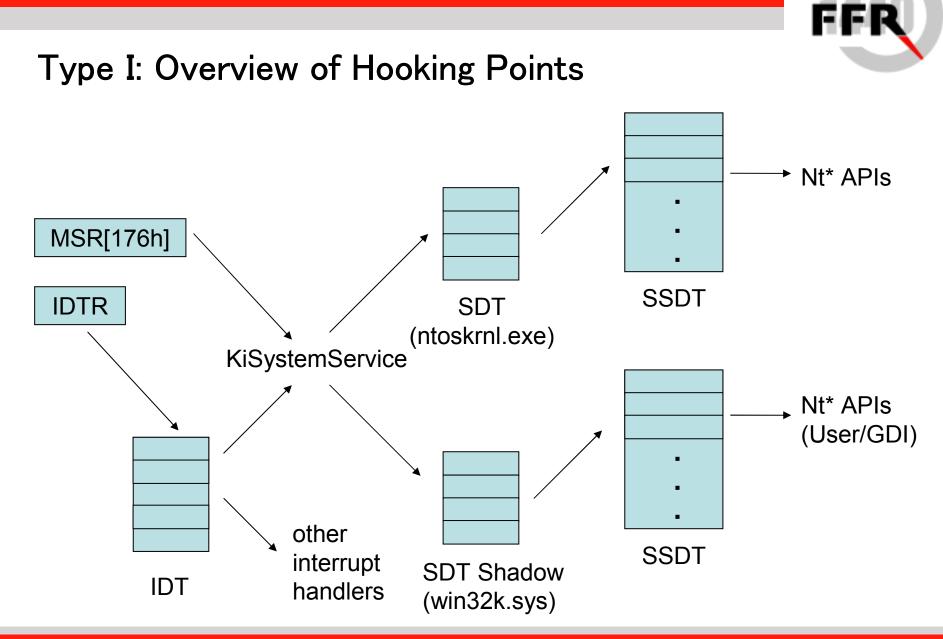


• 1. Review of subversive techniques in kernel space

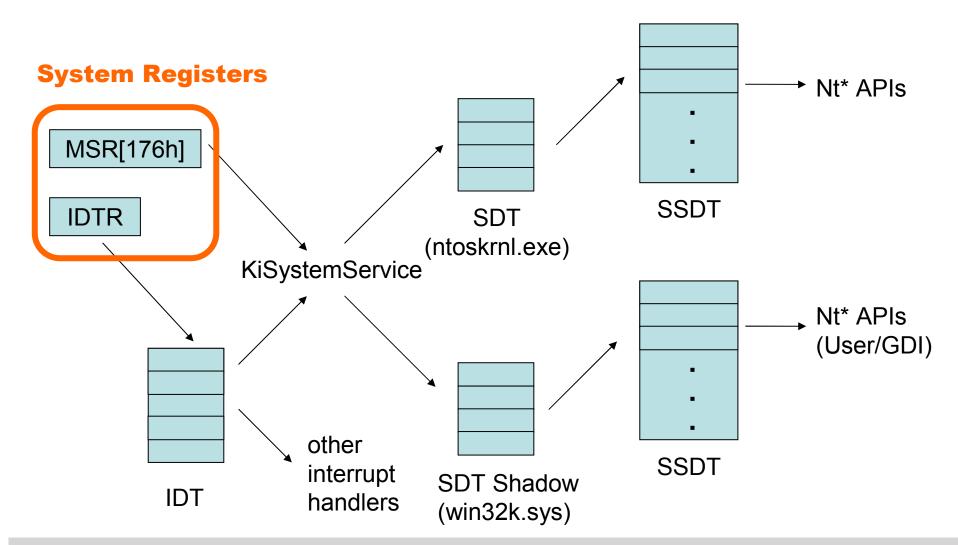


Remember Joanna's classification

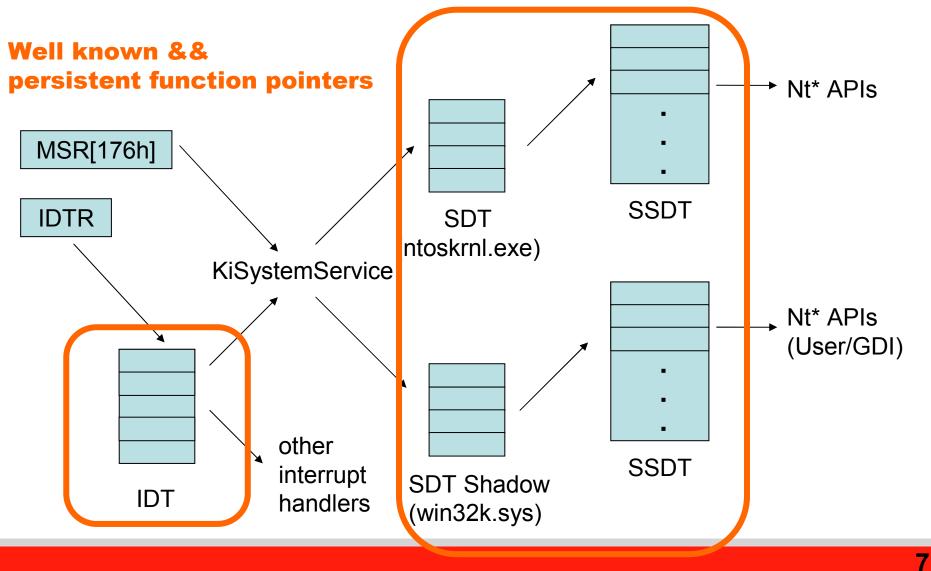
- Joanna Rutkowska proposed stealth malware taxonomy in November, 2006. <u>http://invisiblethings.org/papers/malware-taxonomy.pdf</u>
- Type 0
 - standalone malware, which never changes any system resources
- Type I
 - changes the persistent system resources
- Type II
 - changes the non-persistent system resources
- Type III
 - malware runs outside the system



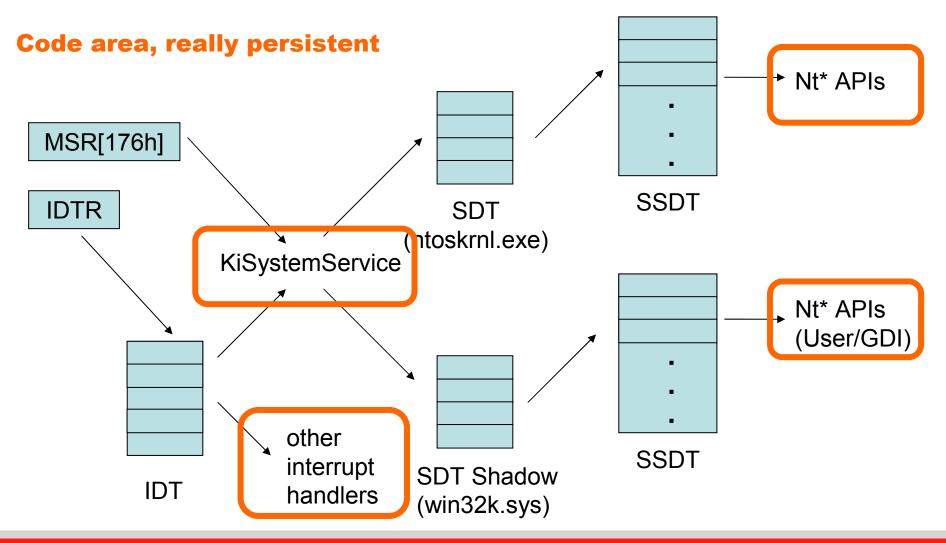
Type I: Overview of Hooking Points



Type I: Overview of Hooking Points



Type I: Overview of Hooking Points





Type I

- It is easy to detect
- PatchGuard in Vista(x64) is a countermeasure for this type
- Many rootkit detectors have been released for this type



Type II

- Malware changes the non-persistent system resources
- · Hooking point might be modified by the regular execution path
- DKOM(Direct Kernel Object Manipulation)
 - by <u>http://www.blackhat.com/presentations/win-usa-04/bh-win-04-butler.pdf</u>
- KOH(Kernel Object Hooking)
 - by Greg Hoglund in Jan, 2006
 <u>http://www.rootkit.com/newsread.php?newsid=501</u>



DKOM(Direct Kernel Object Manipulation)

- Malware manipulates the process list, tokens and other kernel objects directly
- For example:
 - Unlink target process from process list
 - Add/remove priviledges to tokens
- DKOM's possibilities are limited
 - Whether information hiding can be done depends on the implementation of process that deals with the data

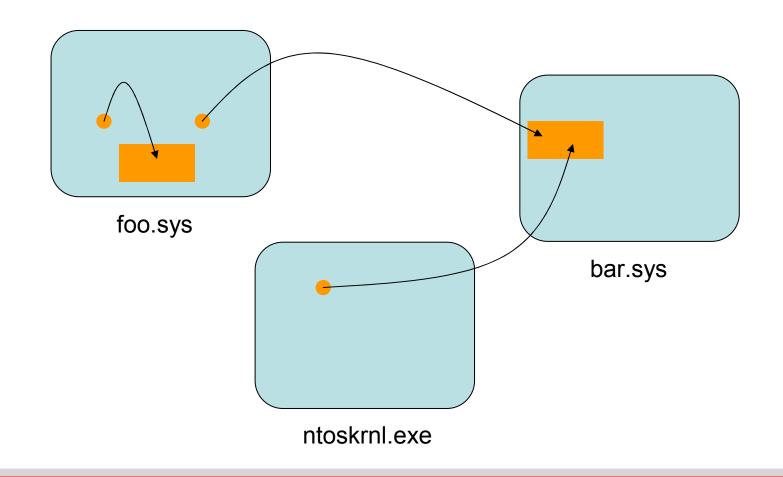


KOH(Kernel Object Hooking)

- Remember the SDT, SSDT and other well known && persistent function pointers?
- Do you know how many such patching points are there in kernel space?
 - They might or might not be persistent
 - It depends on each kernel object
- Detector has to understand all function pointers
- is_within_own_memory_range(PVOID Address) is useful, but not enough



is_within_own_memory_range(PVOID Addr)





Type III

- No malware exists in the system(guest)
- Malware (ab)uses Virtualization Technology
- SMM Rootkit and Firmware Rootkit might also fall into this category (a problem of taxonomy that is not important for our cause)
 - BluePill

.

- Original BP was presented by Joanna Rutkowska in BH-US-2006.
- (Current) New BP supports both Intel VT and AMD-v technologies, and is also capable of on the fly loading and unloading
- BP doesn't modify any system resources on the guest
 - From a technical view, BP patches the guest's PTE to hide its loaded virtual memory from the guest
 - · However this doesn't really help detecting it

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Type III (cont.)

- Vitriol
 - Presented by Dino Dai Zovi, Black Hat US 2006
 - VT-x rootkit, closed source
- VMM Rootkit Framework
 - Posted by Shawn Embleton, Aug, 2007
 <u>http://www.rootkit.com/newsread.php?newsid=758</u>
 - This is really good start point for learning for how to create VMM



Case Study: Storm Worm

- The Storm Worm first appeared in Fall, 2006
- Some variants have rootkit functions to hide from AV products
- As of Jan 2008 we can see "Happy New Year 2008" variants
- When a user clicks onto the executable,



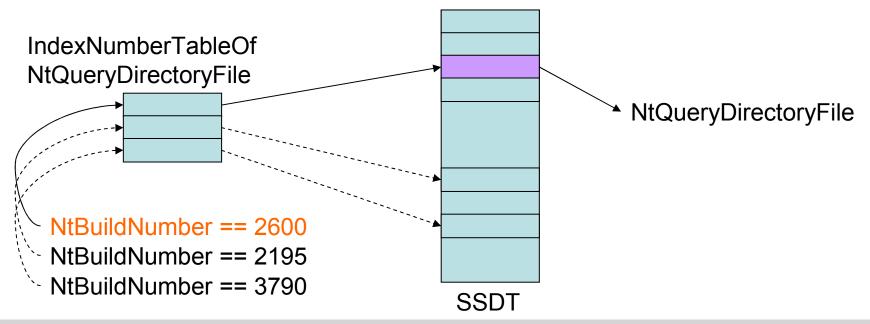
Storm Worm

- 1. Executable drops the system driver (.sys), and loads it into the kernel using Service Control Manager (SCM)
- 2. Driver has two functions shown below
 - Rootkit functions Hide files, registry entries and connections using SSDT and IRP hooking
 - Code Injection function Inject malicious code (not DLL) into process context of services.exe and execute it
- 3. Injected code starts P2P communication



Rootkit functions

- Storm Worm hooks three Native APIs
 - NtQueryDirectoryFile, NtEnumerateKey, NtEnumerateValueKey
- API Index of SSDT is different for each NtBuildNumber
- Storm Worm has index number tables for build 2195(2k), 2600(XP) and 3790(2k3)





Rootkit functions (cont.)

- It hooks the IRP_DEVICE_CONTROL routine by patching the TCP DriverObject's IRP table ("¥¥Device¥¥Tcp")
- Hide connections from netstat

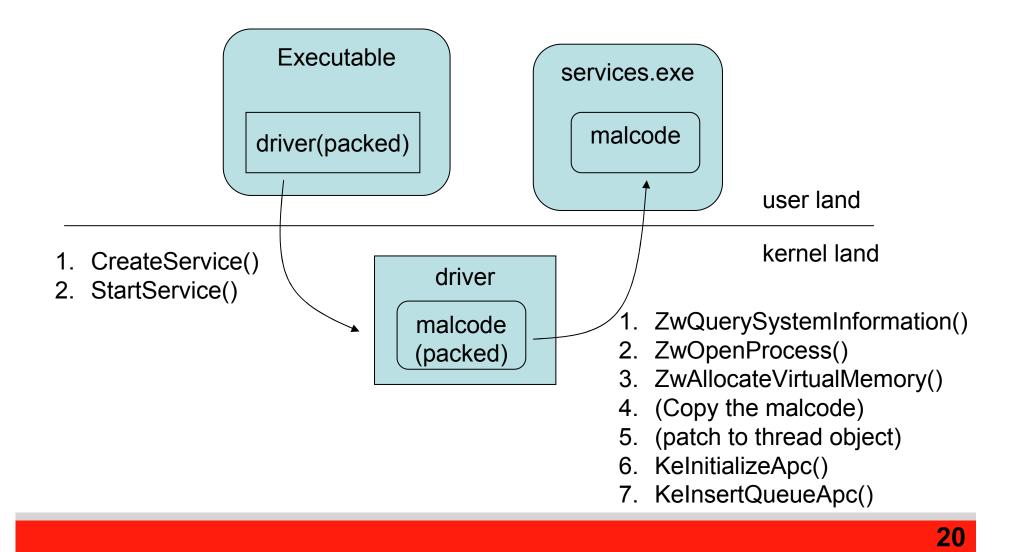
But is this KOH?

YES: It modifies the IRP Table contained within the DriverObject

NO: Many people know about the existence of IRP tables



Code injection function



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2. Review of subversive techniques in kernel space

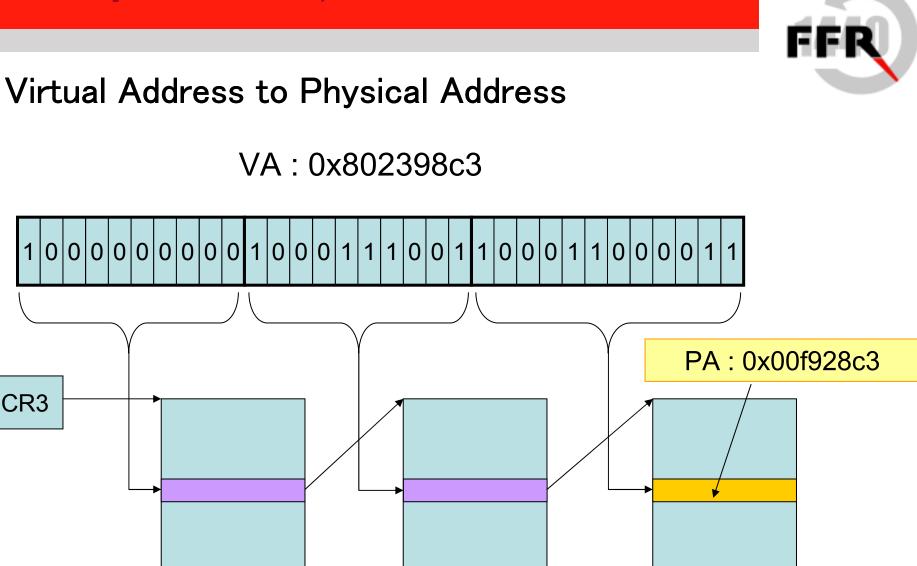


What we have to consider "Virtualization"

- CPU Virtualization
 - Some registers should be reserved for VMM and each VM. GDTR, LDTR, IDTR, CR0-4, DR0-7, MSR, Segment Register, etc
 - Exceptions
- Memory Virtualization
 - should separate VMM memory space and each VM's memory space
- Device Virtualization
 - Interrupt, I/O instructions, MMIO, DMA access

PD

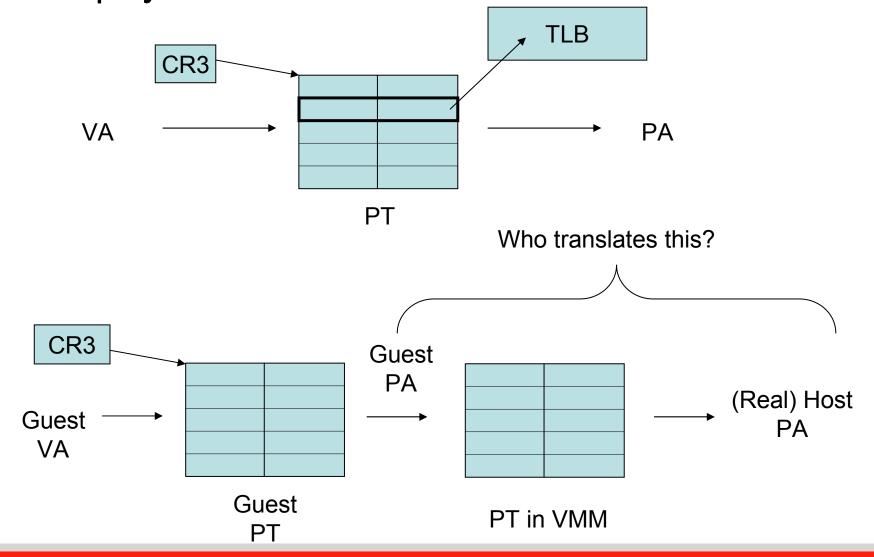
CR3



ΡT

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To simplify...





Memory virtualization

- If the processor supports EPT (Extended Page Table), this 2-stages translation is automatically done by the MMU
 - EPT is not implemented yet
- VMM should implement this translation as software using Shadow Paging

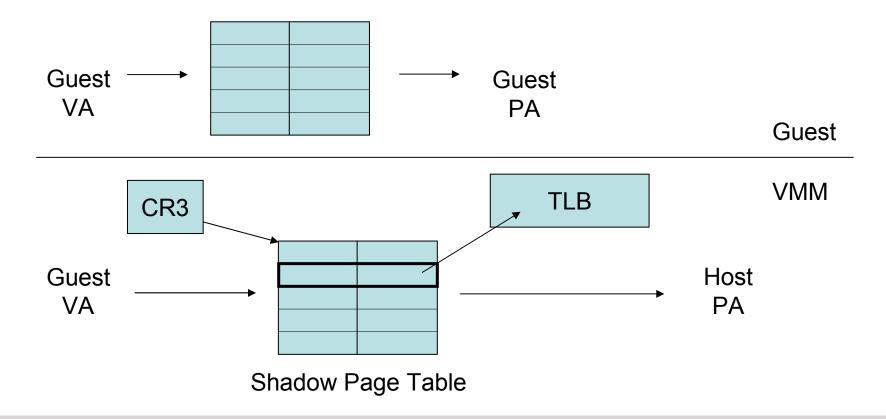
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Shadow Paging

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- VMM updates SPT on #PF in the guest
 - and also emulates TLB flush caused by MOV to CR3 and INVLPG



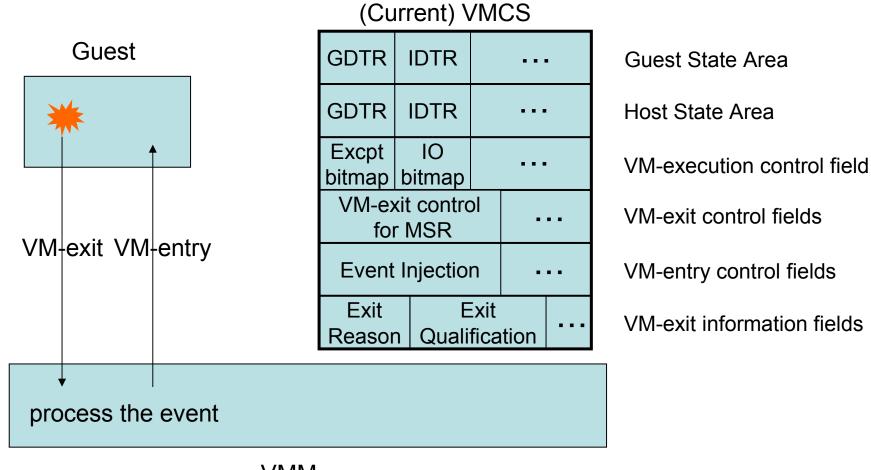


Intel VT

- Intel VT is the Intel VT-* family's generic name
 - VT-x, virtualization for x86/64
 - VT-d, virtualization for device (Directed I/O)
 - VT-i, virtualization for Itanium
- Key factors
 - VMX mode
 - VMX root-operations(ring0-3)
 - VMX non-root-operations(ring0-3)
 - VMCS (Virtual Machine Control Structure)
 - VMX Instructions set
 - VMXON, VMXOFF, VMLAUNCH, VMRESUME, VMCALL, VMWRITE, VMREAD, VMCLEAR, VMPTRLD, VMPTRST

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How Intel VT works:



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enum EXIT_REASON {

- Specific instructions
 - CPUID, INVD, INVLPG, RDTSC, RDPMC, HLT, etc.
 - All VMX Instructions
- I/O Instructions
 - IN, OUT, etc.
- Exceptions
- Access to CR0-CR4, DR0-DR7, MSR
- etc.

};



Steps to launch the VMM and VM

- Confirm that the processor supports VMX operations
 - CPUID
- · Confirm that VMX operations are not disabled in the BIOS
 - MSR_IA32_FEATURE_CONTROL
- Set the CR4.VMXE bit
- Allocate and Initialize VMXON region
 - Write lower 32 bits value of VMX_BASIC_MSR to VMXON region
- Execute VMXON
 - CR0.PE, CR0.PG, and CR4.VME must be set.



Steps to launch the VM and VMM (cont.)

- Allocate VMCS regions
- Execute VMPTRLD to set Current VMCS
- Initialize Current VMCS using VMREAD and VMWRITE
 - VMCS contains the EP of VMM, and Guest IP after VMLAUNCH
- Execute VMLAUNCH
 - Continue to execute the guest from IP is contained in VMCS
- When VM-exit occurred, IP and other registers are switched to VMM ones.

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3. Viton, Hypervisor IPS

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Viton

- IPS, which runs outside the guest
- Just a PoC, tested on Windows XP SP2 only
- Force immutability to persistent system resources
- Observe control/system registers modification, and VMX instructions are raised in the guest
- Offer the extensibility for monitoring the guest activity

• It is based on **Bitvisor**

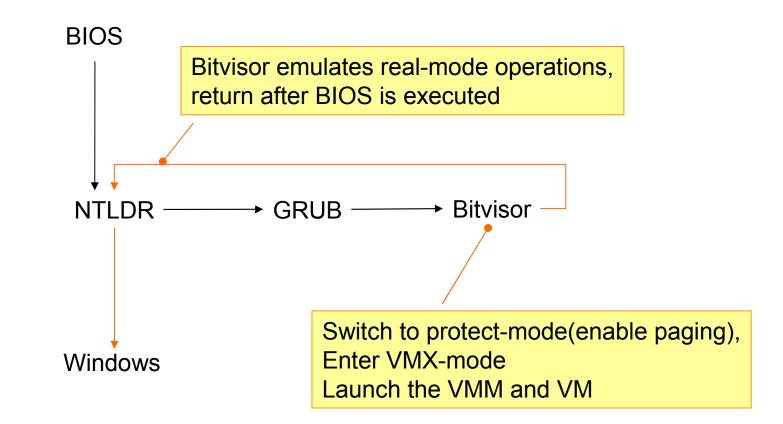


Bitvisor - <u>http://www.securevm.org</u>

- The Bitvisor VMM software is developed by the Secure VM project centered around Tsukuba Univ. in Japan
- Features:
 - Open source, BSD License
 - Semi-path through model
 - Type I VMM (Hypervisor model, like Xen)
 - Full scratched, pure domestic production
 - Support for 32/64 bits architecture in VMM
 - Support for Multi-core/processor in VMM and Guest
 - Can run Windows XP/Vista as Guests without modification
 - Support for PAE in the Guest
 - Support for Real-mode emulation



How Bitvisor works: Launch process



What Viton protects/detects:

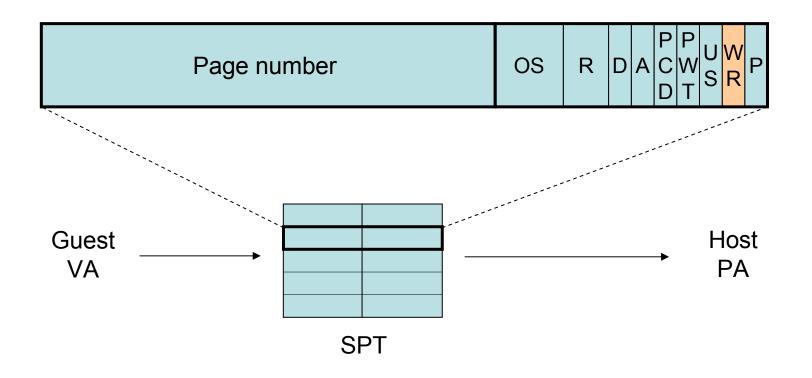
- Instructions
 - Detect and block all VMX Instructions
- Registers
 - Watchdog for IDTR
 - Locking the MSR[SYSTENR_EIP]
 - Locking the CR0.WP Bit
- Memory
 - Protect from modification
 - All code sections (R-X) in ntoskrnl.exe
 - · IDT
 - · SDT
 - SDT.ST (SSDT)



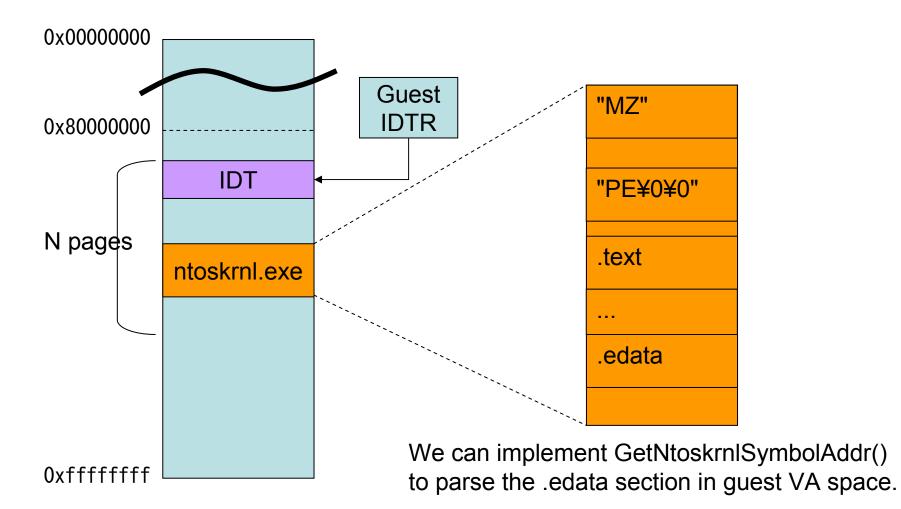


How to protect the guest memory modification

- Viton clears the WR bit in a SPT entry
 - If CR0.WP is set, even the kernel cannot modify the page



How to recognize the guest memory layout



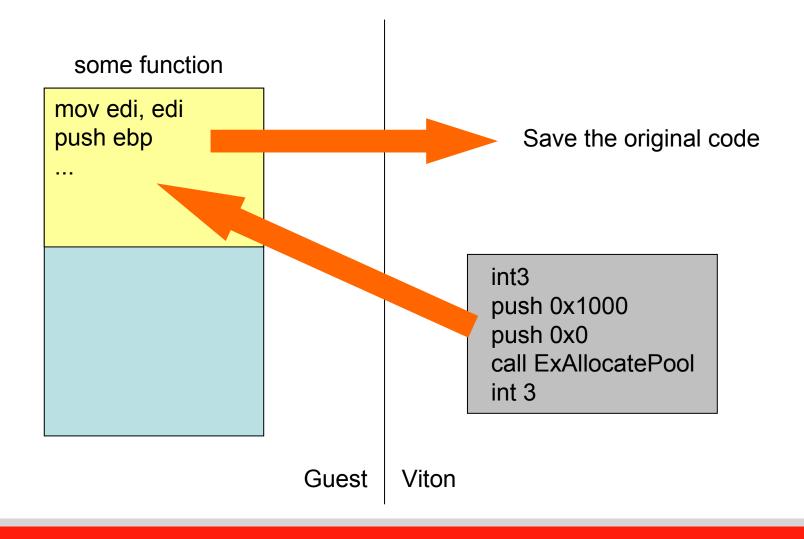


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Guest activity monitoring

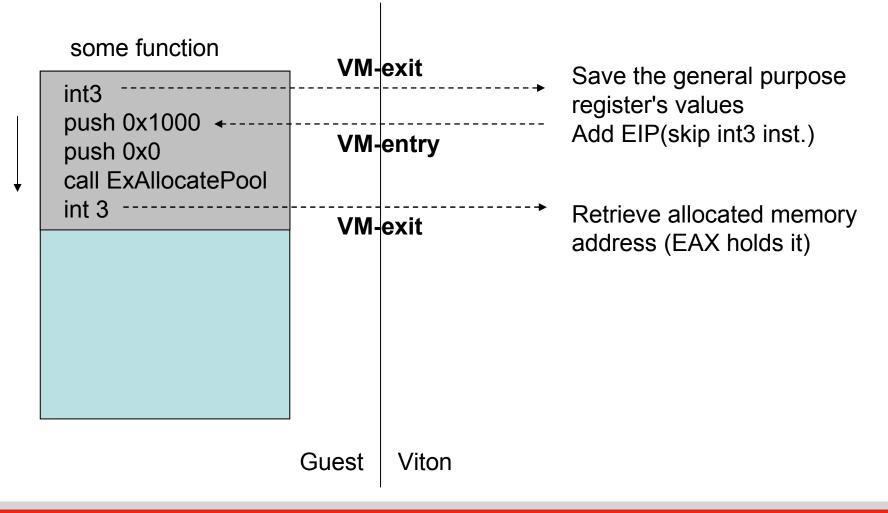
- When we use the Viton, no one can modify the kernel code, excluding the Viton.
- Viton can monitor the guest's activity by hooking the code
 - 1. Allocate memory for detours in the guest VA space
 - 2. Setup the detours buffer
 - 3. Hook the target function

How to allocate memory in guest VA space



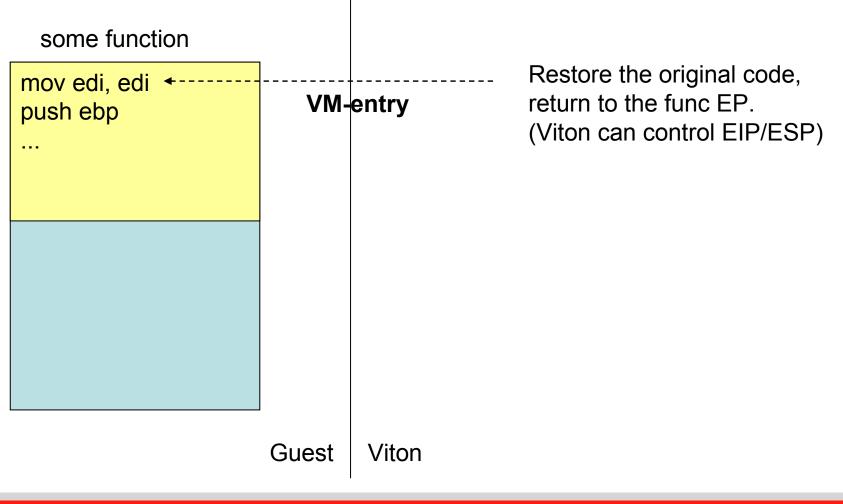


How to allocate memory in guest VA space





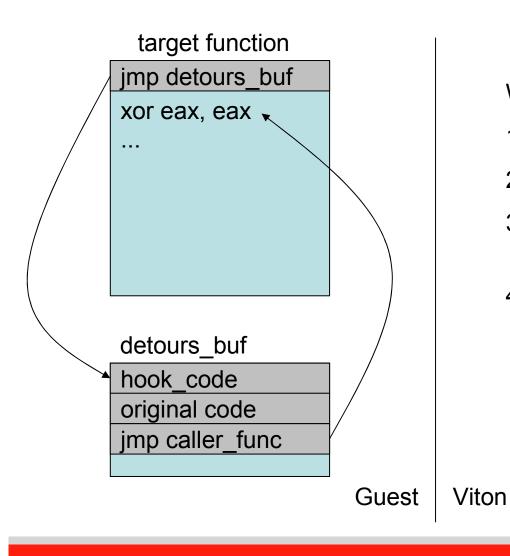
How to allocate memory in guest VA space



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How to hook the guest code



When the target function is called,

- 1. jump to the detours_buf
- 2. Execute our hook_code
- 3. Execute original code which is overwritten by "jmp detours_buf"
- 4. jump to the next code of overwritten one



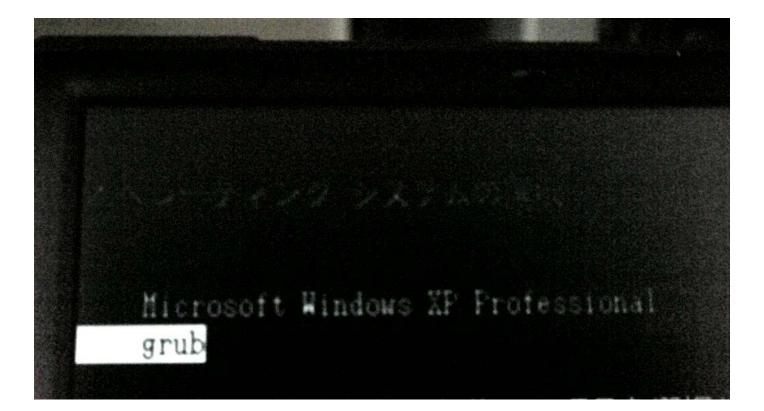
What can Viton do hooking the guest code ?

- Viton can retrieve the guest information in hook_code
 - int3 and other inst. that cause VM-exit are useful
- So, Wouldn't you hook below functions ?
 - ZwCreateProcess/ZwTerminateProcess
 - ZwLoadDriver
- Then, Viton understands process, driver and other guest system resource information.



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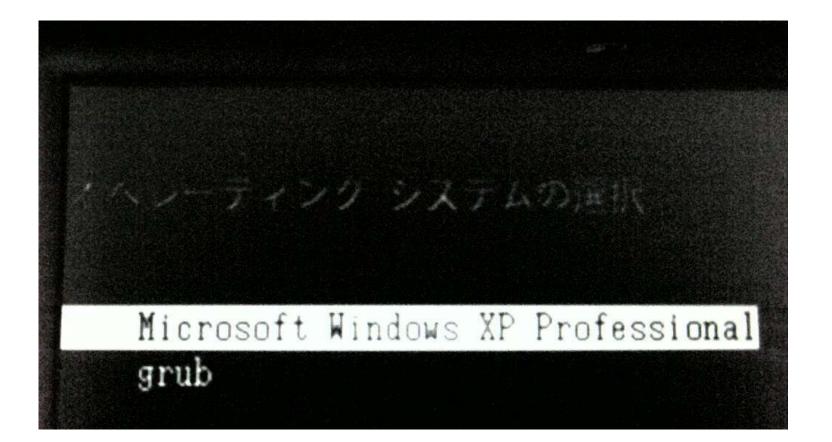
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Viton		
	and load	NTLDR of Windows NT/2K/XP
		CMLDR of Windows NT/2K/XP
		10.SYS of Windows 9x/Me
		Mandriva with menu. 1st alre
		Linux with menu. 1st already
	dline	
	(fd8)	
	to dos	
reboot		
halt		

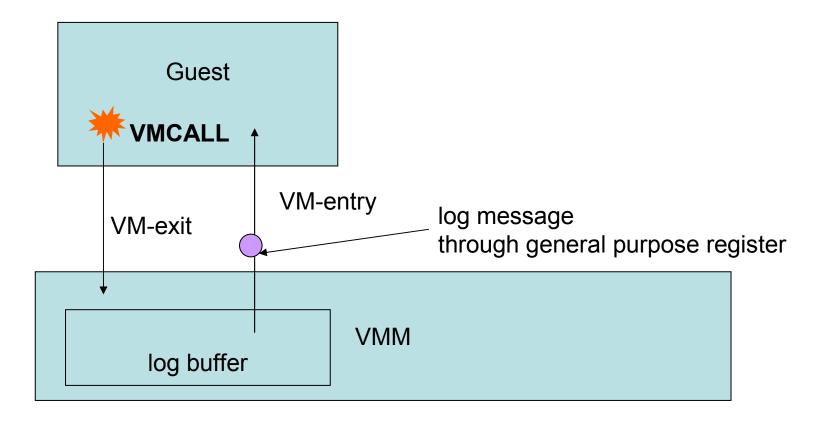
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dbgsh (Bitvisor's debuging function)



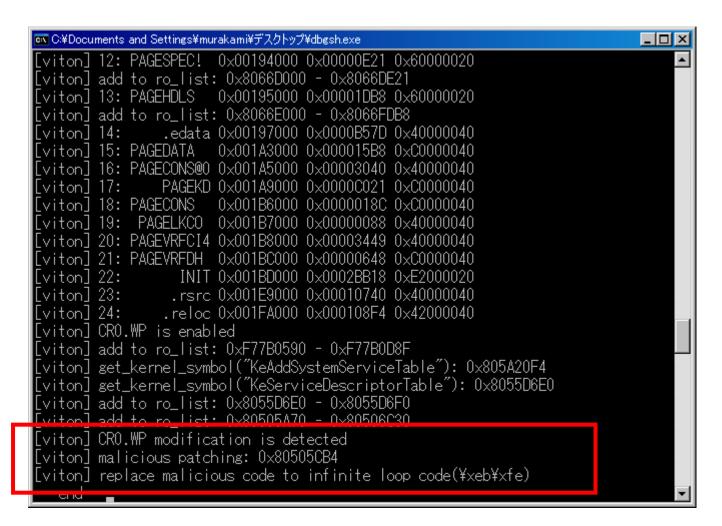
🗪 C#Documents and Settings¥murakami¥デスクトップ¥dbgsh.exe	- D ×
[viton] CR0.WP is enabled	_
[viton] VMCS_GUEST_IDTR_BASE: 0x8003F400	
[viton] Guest IDT[0x2E].handler: 0x80542000	
[viton] Start address for searching kernel base: 0x80442000	
[viton] Kernel base: 0x804D9000	
[viton] 0: .text 0x00001000 0x0006E46C 0x68000020	
[viton] add to ro_list: 0x804DA168 - 0x8054846C	
[viton] 1: POOLMI 0x00070000 0x000011F9 0x68000020	
[viton] add to ro_list: 0x80549000 - 0x8054A1F9	
[viton] 2: MISYSPTE 0x00072000 0x000006CB 0x68000020	
[viton] add to ro_list: 0x8054B000 - 0x8054B6CB	
[viton] 3: POOLCODE 0x00073000 0x000012AE 0x68000020	
[viton] add to ro_list: 0x8054C000 - 0x8054D2AE	
[viton] 4:	
[viton] 5: INITDATA8 0x0008E000 0x00000038 0xC8000040	
[viton] 6: INITCONSe 0x0008F000 0x00001A65 0x48000040	
[viton] 7: PAGE 0x00091000 0x000DECDF 0x60000020	
[viton] add to ro_list: 0x8056A000 - 0x80648CDF	
[viton] 8: PAGELK 0x00170000 0x0000E520 0x60000020	
[viton] add to ro_list: 0x80649000 - 0x80657520	
[viton] 9: PAGEVRFY	
[viton] add to ro_list: 0x80658000 - 0x80666AA6 [viton] 10: PAGEWMI 0x0018E000 0x00001703 0x60000020	
[viton] 10: PAGEWMI 0x0018E000 0x00001703 0x60000020 [viton] add to ro_list: 0x80667000 - 0x80668703	
	-
[viton] 11: PAGEKD 0x00190000 0x00003D93 0x60000020	

ro_list: read only list

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OSR Driver Loader			<u>? ×</u>
Open System: 105 Route 10 Amherst, NH Ph: (603) 595 Fax: (603) 595 Ver: V2.3 - 1	03031 -6500 5-6503		Exit Help ServiceGroupOrder Active Services
Registry Key:	stormworm_driver		
Driver Path:	nd Settings¥muraka	ami¥デスクトッ <mark>が</mark> ¥stormworn	n_driver.sys <mark>B</mark> rowse
Driver Version: Driver Size: Driver File Time:	129536 Bytes Monday, January 07,	, 2008 16:07:00	
Display Name:	stormworm_driver		
Service Start:	Demand		•
Load Group: Order In Group: Depend On Group(s):	None Type: AudioGroup Base	Driver Er	<u>G</u> roup Load Order ror: Normal ▼
	Boot Bus Extender Boot File System		
Last Status:	この操作を正しく終了し	しました。	
MiniFilter Settings Default Instance: Altitude AndFlags		Altitude: 0 Flags: 0	
<u>R</u> egister Service	<u>U</u> nregister Service	Start Service	S <u>t</u> op Service







OSR Driver Loader			<u>? ×</u>
Open System: 105 Route 10 Amherst, NH Ph: (603) 595 Fax: (603) 595 Ver: V2.3 -	03031 6500 56503		<u>E</u> xit <u>H</u> elp ServiceGroup <u>O</u> rder <u>A</u> ctive Services
Registry Key: Driver Path:	vmxcpu0 Juments and Settings	s¥murakami¥デスクトッ	⁷ ¥vmxcpu0.sys Browse
Driver Version: Driver Size: Driver File Time:	22400 Bytes Thursday, August 07	7, 2008 20:48:06	
Display Name: Service Start:	vmxcpu0 Demand	1	
Load Group: Order In Group: Depend On Group(s):	None Type:	Driver 💌	Group Load Order Error: Normal
Depend on Group'ss.	AudioGroup Base Boot Bus Extender Boot File System		▲ ▼
Last Status:	システムに接続されたう	デバイスが機能していま	せん。
MiniFilter Settings Default Instance: AltitudeAndFlags		Altitude: 0 Flags: 0	
<u>R</u> egister Service	<u>U</u> nregister Service	Start Service	Stop Service



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Eile		ure Options Co <u>m</u> puter <u>H</u> elp	
#	Time	Debug Print	
)	7:08:58	[vmm] Driver Routines [00000000]	
	7:08:58	[vmm] [0000000]	
2	7:08:58	[vmm] Driver Entry [A605F4A0]	
3	7:08:58	[vmm] Driver Unload [A605F3B0]	
1	7:08:58	[vmm] StartVMX [A605D640]	
5	7:08:58	[vmm] VMMEntryPoint [A605F830]	
ì	7:08:58	[vmm] VMXONRegion virtual address [A29B1000]	
7	7:08:58	[vmm] VMXONRegion physical address [20405000]	
3	7:08:58	[vmm] VMCSRegion virtual address [A29AD000]	
9	7:08:58	[vmm] VMCSRegion physical address [185C6000]	
0	7:08:58	[vmm] FakeStack [8475B000]	
1	7:08:58	[vmm] Guest Return EIP [A605F76E]	
2	7:08:58	[vmm] Enabling VMX mode on CPU 0 [00000000]	
3	7:08:58	[vmm] Running on Processor [00000000]	
4 5	7:08:58 7:08:58	[vmm] GDT Base [8003F000] [vmm] GDT Limit [000003FF]	
5 6	7:08:58	[vmm] IDT Base [8003F400]	
7	7:08:58	[vmm] IDT Limit [000007FF]	
8	7:08:58	[vmm] VMX Support Not Present. [0000E39D]	
9	7:08:58	[vmm] Bunning on Processor [00000000]	
20	7:08:58	[vmm] ERROR : Launch aborted. [00000000]	



Viton vs.

- Type I
 - Easy
- Type II
 - DKOM: Difficult, but possible
 - KOH: Difficult, we need more research, and breakthrough
- Type III
 - Easy (First come, first served)



4. Conclusions

- · Virtualization Technology becomes a help to protect the kernel
- However, it is not a silver bullet
 - Foundation for existing security solutions



Thank you!



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