more embedded systems

FX of Phenoelit

Las Vegas 2003

Agenda

- Hacking the Matrix
- GSM 3G Basics
- GPRS backbone hacks
- Anonymous HTTP via WAP
- Siemens S55 Vulnerabilities
- Phenoelit's usual Cisco Øday vulnerability and exploit

Hacking the MATRIX

- Enterasys Matrix E1
- Vulnerabilities in the Matrix:
 - SSH can only fork 10 times
 - unfinished connections stay open
 - TCP ISN 64k rule on switch ports
 - OSPF neighbors added to neighbor list in state HELO
 - HTTP Server negative content-length integer bug

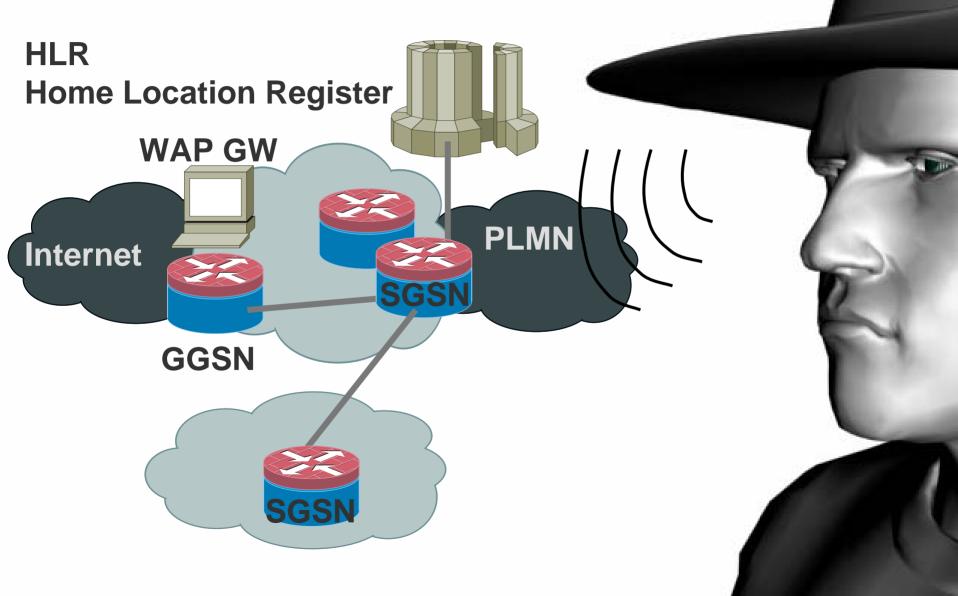
GSM Basics

- Global System for Mobile communications – Mobile Phone Network
- Authentication based on key material on SIM card and on the network
- GSM core network relies on caller identification via MSISDN and key material for authentication and billing
- Caller spoofing generally considered hard to do

GSM Basics – GPRS

- General Packet
 Radio Service
- Packet oriented data transfer for mobile devices
- Backbone build using TCP/IP
- Authentication via GSM and optional PPP (CHAP,PAP)
- Network access, routes and IP addressing depending on Access Point Name (APN)

GPRS Components



GPRS Attack Points

- The GGSN is just another
 TCP/IP device facing the Internet
 - @stake: Nokia GPRS 1 IPSO DoS with TCP Option 0xFF
- APN guessing (WarAPNing?)
 - APNs often selected by company name for mobile VPNs
 - APN filtering in HLR possible but rarely implemented
 - APNs are not considered a secret ③

Backbone hacks: GTP

- GPRS Tunneling Protocol (ETSI TS 129 060)
- Transports user protocols in the GPRS backbone
- Controls inter-network roaming
- Control channel and user data channel
- xGSN has to support all old protocol versions from 0 on

What's so cool?

- Fully UDP based protocol
- No authentication
- Full control over all GPRS nodes
- Creation, update or deletion of user contexts and tunnels
- Rerouting, redirection and relocation control
- Context billing configuration
- "Invitation" of mobile users to APNs
- Forced roaming of users

What's so cool?

ETSI TS 129 060 / Section 12: No security is provided in GTP to protect the communication between different GPRS networks. [...] A security mechanism that may be considered is for example IP Security.

Anonymous HTTP

- Wireless Application Protocol uses
 - Wireless Transport Protocol (WTP)
 - Wireless Session Protocol (WSP)
- WSP uses connections over UDP
 - Not easily spoofable due to 32bit session IDs
- WSP supports connectionless transactions
 - No acknowledge messages
 - Single UDP packet transfer
 - Full HTTP Request capabilities

Anonymous HTTP [2]

- 1. GPRS WAP APN or open WAP Gateway
- 2. Send HTTP request to WAP Gateway using someone else's IP address
- 3. Destination UDP port 9200
- 4. Enjoy

Note: Victim IP addresses can be collected by running a WAP site using HTTP header field "X-Forwarded-For".

Siemens S55 – Bluetooth

- Pairing to death
 - Every connection creates a dialog
 - Connection structures not cleared
 - Bluetooth connection exhaustion
 - All you need is l2ping
- The Big Inbox
 - S55 accepts any file sent to it
 - Sending 2100 files is possible
 - Deleting 2100 files is not

Siemens S55 – Java

- NULL Pointer Exceptions
 - Generally crash JVM (white screen of death)
- JAD file parsing
 - Missing "MIDlet-Name" and "MIDlet-1" tags crash JVM
 - Overlong "MIDlet-Name" leads to buffer overflow in filesystem



Siemens S55 – Java

MIDlet-Name: Test

File name:

/Java/jam/\${MIDlet-Name}/thing.jar

 \rightarrow /Java/jam/Test/thing.jar

MIDlet-Name: AAA...AAA
File name:
/Java/jam/\${MIDlet-Name}/thing.jar
→ /Java/jam/AAA...AAA/thingAAA.jar

Siemens S55 – Spy.jar

- Outgoing WAP connections
 - Java Applications can connect to WAP sites without user permission
 - Com.siemens.mp.gsm.
 PhoneBook.getMDN() allows access to missed call list
 - System.getProperty("IMEI") returns the International Mobile Equipment Identity

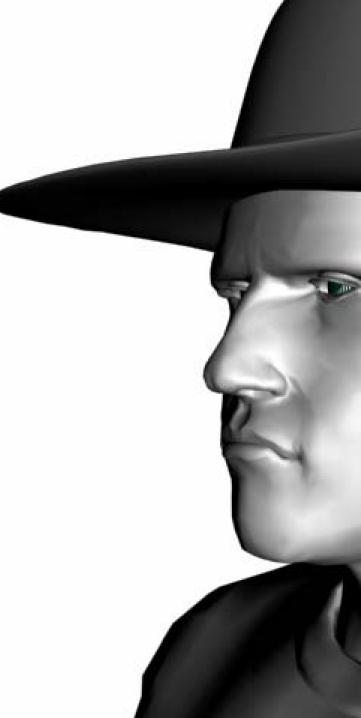
Siemens S55 – Time.jar

- Sending SMS or placing calls via Java applications requires user permission
- Permission is obtained via dialog
- Filling the screen obscures the dialog
- User answers a different question
- Outgoing call triggered but terminates Java application

Siemens S55 GIFt

- S55 supports GIF for MMS
- GIF file format has a Virtual Screen section
- Changing the virtual screen offset for a picture or one frame in an animated GIF crashes the device
- Placing such a pic as background renders the device unusable

And Now For Something Completely Different



A small bug ...

- Cisco IOS 11.x and below
- UDP Echo service memory leak
 - Device sends as much data back to the sender as the UDP length field said it got
 - Leaks IO memory blocks
 - IO memory contains actual packet data – and not just ours
 - We are talking about 19kbytes here

IOS Fingerprinting

- Leaked IO memory contains memory block headers
 - Block headers contain address of who allocated the block
 - Address of allocating function changes per image
 - Address range changes per platform
- Result:

Reliable remote IOS fingerprint

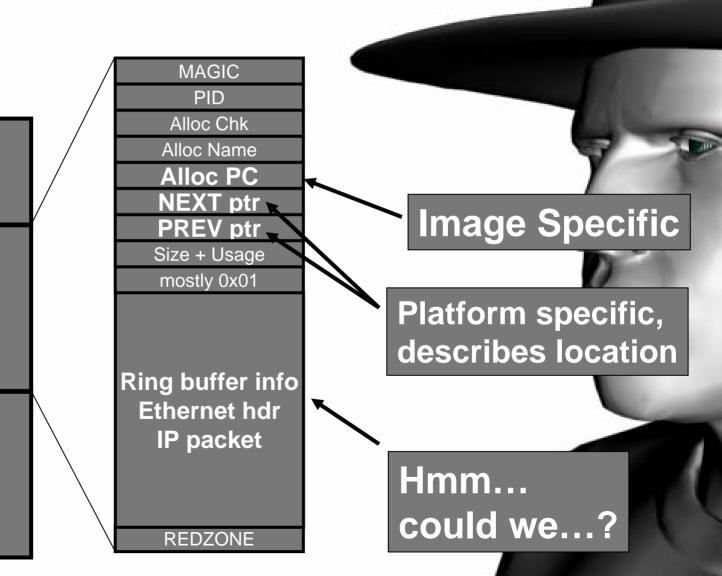
IOS Fingerprinting [2]

In detail:

Echo Data 0x00 ...

Receive Buffer

Receive Buffer



Remote IOS Sniffing

- Leaked IO memory contains packets in the receive buffers (RX ring ds elements)
- Phenoelit IOSniff
 - Repeated memory leak retrieval
 - Memory block identification
 - Packet offset identification
 - Packet decoding
 - Caching and duplicate prevention

Remote IOS Sniffing

```
[0x00E0B42C]: 00:60:47:4F:5E:72 -> 01:00:0C:CC:CC:CC
pure Ethernet stuff
 .....+...radio.b.phenoelit.de.................Ether
net0.....Cisco Internetwork Operating System Software
  .IOS (tm) 1600 Software (C1600-Y-L), Version 11.3(11b), REL
EASE SOFTWARE (fc1).Copyright (c) 1986-2001 by cisco Systems
 , Inc..Compiled Fri 02-Mar-01 17:12 by cmong....cisco 1603..
[0x00E0CF2C]: 00:A0:24:2B:BE:BB -> 00:00:0C:4A:9C:C2
192.168.1.3 -> 192.168.1.16 43 bytes [TTL 63] DF (payload 23)
[TCP] 1035 -> 23 (783944042/983338029) ACK PSH win 32120
(payload 3)
en.
[0x00E112AC]: 00:A0:24:2B:BE:BB -> 00:00:0C:4A:9C:C2
192.168.1.3 -> 192.168.1.16 46 bytes [TTL 63] DF (payload 26)
[TCP] 1035 -> 23 (783944045/983338043) ACK PSH win 32120
(payload 6)
s3cr3t.
[0x00E1196C]: 00:00:0C:4A:9C:C2 -> 00:01:03:8C:9B:44
[ARP] Reply for 192.168.1.100 from 192.168.1.16 (MAC:
00:01:03:8C:9B:44)
```

IOS HTTP bug

- Almost all embedded HTTP implementations are vulnerable – Cisco is no exception
- Integer or counting related issue
- IOS 11.x 12.x
- Requires sending of a 2GB sized URL to the device
- Stack based buffer overflow

Exploitation – issues in the past

- For heap overflows, we need several image and configuration depend addresses
 - PREV pointer in the memory block
 - Size value in IO memory exploitation
 - Stack location
 - Own code location
- Requirements made reliable remote exploitation hard / impossible

What we got now

- UDP Echo memory leak
 - Attacker provided binary data (the delivered Echo content)
 - Live IOS memory addresses (leaked IO memory block headers)
 - Ability to fill multiple memory areas with our binary data (Ring buffer)
- HTTP Overflow
 - Direct frame pointer and return address overwrite

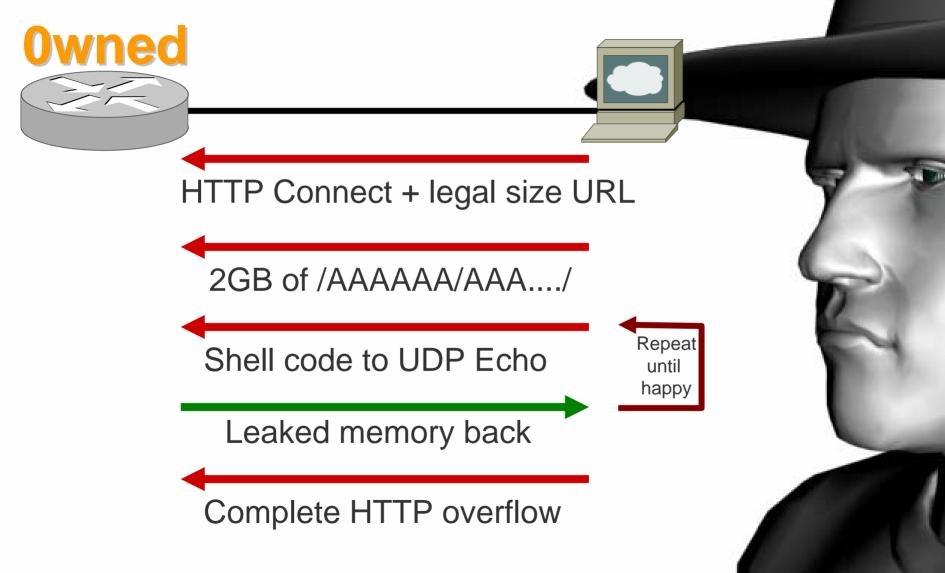
What we can do now

- Send full binary shell code
- Calculate the address of the code using IO memory block header information
- Select the shell code that is most likely not modified
- Directly redirect execution in the provided shell code
- Own the box

Combining

- 1. Send the maximum URL length allowed by IOS
- 2. Send 2GB of additional URL elements in correct sized chunks
- 3. Perform UDP memory leak several times with shell code in the request packet
- 4. Make intelligent decision on which address to use
- 5. Complete overflow and gain control

Again, in color



Binary via HTTP

- Cisco's HTTP doesn't like all characters
 - Slash, 0x0a, 0x0d and 0x00 are obviously bad for HTTP
 - Some others are bad as well
- HTTP encoding (%XY) supported
- Decoding seems to take place in the exact same buffer
- Return address HTTP encoded

Return address selection

- Several address selection strategies tested
 - Last address obtained (about 50% success)
 - Randomly selected address (about 50%-60% success)
 - Highest memory location (about 0%-10% success)
 - Lowest memory location (about 90% success)
 - Most frequently seen address (about 30%-40% success)

Cisco shell code in the past

- Complete configuration replacement in NVRAM and requires reboot
- Required knowledge
 - Attacked interface
 - IP address
 - Basic routing information
- Looses information of original configuration
 - Passwords and keys
 - Other routing information
 - Access lists
 - Logging information

Researching binary IOS

- Cisco supports serial gdb
- ROM Monitor (rommon) allows limited debugging
 - Breakpoints
 - Watchpoints
 - Disassembly
- Code identification simple
 - Related debug strings can be found in the code
 - Data and text segment are intermixed with each other
 - Strings stored before the function

Next generation code

- Runtime IOS patching things of interest:
 - Disable IOS text segment checksum function
 - Patch out authentication requirement for incoming VTY connections
 - Patch out verification return code from "enable mode" function
 - ACLs or BGP neighbor check?
- Keep IOS running ... but how?

Clean return

- Overflow destroys significant amounts of stack due to HTTP encoding
 - 24 bytes encoded: %fe%fe%ba%be%f0%0d%ca%fe
 - 8 bytes decoded
- Motorola call structure uses frame pointer in A6 and saved stack pointer on stack
- Moving the stack pointer before the saved SP of any function restores SP and A6
- Search stack "upward" for return address of desired function
 - SP = <current> 4
 unlk a6
 rts

Clean return code

IOS 11.3(11b) HTTP overflow find-return code

move.1

findret: addq.1 #0x01,a7 cmp.1 #0x0219fcc0,(a7) findret bne move.l a7,(a2) sub.l #0x0000004,(a2) move.l (a2),a6 clr.l **d**0 -4(a6),a2movem.1 unlk **a**6 rts

a7,a2

Runtime IOS patching

- Advantages
 - Router stays online
 - Configuration preserved
 - Backdoor in IOS runtime code
- Disadvantages
 - Depending on image
 - Large target list required (code addresses per image)
 - Annoying "checksum error" message on console ☺

Phenoelit HTTP exploit

- Reliable remote IOS exploitation
- Address calculation and shell code placement via UDP Echo info leak
- Address selection using second smallest address
 - first used for HTTP transfer itself
- Runtime IOS patch disables VTY and enable mode password verification

Phenoelit HTTP exploit

Trying... Connected to c1600.mgmt.nsa.gov. Escape character is '^]'.

radio>en
Password:
Password:
Password:
% Bad secrets

radio#sh ru Building configuration...

So what?

- Didn't we mention, that ...
 - ... you shouldn't run unneeded services
 - ... you should protect your infrastructure
 - ... you should not copy data into buffers that are not large enough to hold it
- IOS moves "forward"
- Legal interception is build into IOS "My other computer is your legal interception system." – Hey Jaya, it's no longer a joke.
- If your infrastructure is owned, you can't defend your systems.
- Other people exploit IOS as well only we do it in public.

So what?

"Body of Secrets", James Bamford: By looking for vulnerabilities in Cisco Routers, the NSA can find and capture a lot of electronic messages.

NSA Director Terry Thompson: "But today, I really need someone who knows Cisco routers inside-out and helps me understand how they are used in target networks."

Defense – Mobile Phones

- Turn off unneeded interfaces on your phone
 - Bluetooth
 - IRDA
- Do not run Java code you don't trust
- When receiving files, delete them directly instead of opening them (remember your Outlook)
- Keep your Phone firmware up to date
- Don't use GPRS based VPN solutions

Defense - Cisco

- Do not trust devices, just
 because they are in a black box
- Keep your IOS up to date
- If possible, block direct communication to your infrastructure devices
- Do not run unneeded services on routers
- Prefer out-of-band management
- Include your router in your IDS watch list

Thanks and Greets

- Phenoelit
- Jeff & Ping
- Halvar
- Gaus at Cisco
- All the people being at ph-neutral
- You, the Black Hat audience

