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Taking DMA Attacks to the Next Level:

How to do arbitrary reads/writes in a live and unmodified system using a rogue memory controller

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Who are we

- Anna Trikalinou, PhD Research Scientist Intel Labs, Intel Corporation Email: anna.trikalinou@intel.com
- Dan Lake

Systems Engineer Intel Labs, Intel Corporation Email: dan.lake@intel.com



About this talk

- This talk describes early research into a <u>physical attack</u> on DRAM memory
- We describe <u>hardware design</u> and <u>protocol limitations</u>
- If successful, the result is <u>full access to all physical memory</u>
- Applicable to all memories which follow the JEDEC spec and all architectures



Intro to DMA attacks

- Physical attack where the attacker connects to a DMA-capable port and gains full read/write access to the entire physical memory
 - e.g. PCI, PCI-E, FireWire, USB, etc.
- Goal:
 - Extract secrets (i.e. disk encryption keys)
 - Bypass platform's security policies (i.e. modify page tables)



Example of DMA attack

- Ulf Frisk demonstrated a DMA attack on a MacBook Air
- He was able to obtain the FileVault disk encryption password
- macOS Sierra 10.12.2 software update patched the security issue (Dec 2016)



Source: http://blog.frizk.net/2016/12/filevault-password-retrieval.html



Pros & Cons of Existing Attacks from Attacker's Perspective

- Pros:
 - HW, SW and tutorials are readily available
- Cons:
 - Require specific interface
 - Can be mitigated by BIOS & VT-d
 - Can be mitigated by blocking associated drivers and ports



Motivation

- Pros:
 - HW, SW and tutorials are readily available
- Cons:
 - Require specific interface
 - Can be mitigated by BIOS & VT-d
 - Can be mitigated by blocking associated drivers and ports

What if we could eliminate these?



DRAM design

- A DIMM consists of a series of DRAM chips, mounted on a PCB.
- The DIMM is inserted into a DIMM socket on the motherboard, which then communicates with the processor.





Photo credits: https://www.oempcworld.com/mm5/graphics/00000001/ddr4.jpg http://p.globalsources.com/IMAGES/PDT/B1123866513/DIMM-Socket-Connector.jpg 9



DRAM Placement

4 available DIMM Sockets 1 DIMM inserted

Reference System: Desktop with current generation motherboard and CPU













DRAM Placement on Back Side

Reference System: Desktop with current generation motherboard and CPU





DRAM HW Design

- <u>All 288-pins</u> of each DIMM socket <u>are exposed</u> on both sides of the motherboard
- Pins are electrically connected





How could we exploit that?

• What if we could plug into those pins and request reads/writes from the DIMM as a "memory controller"?



DRAM Functionality



- Specified by JEDEC Standard
- Defines the set of requirements that must be satisfied by all memory modules and all architectures





- 1. CPU reads DIMM's Serial Presence Detect (SPD) EEPROM data via SMBus
 - Manufacturer, module serial, supported voltage, minimum latencies, etc.
 - Bypass Memory Controller







2. CPU decides on a clock frequency and configures Memory Controller





- 3. Memory Controller performs memory calibration
 - MC is agnostic of motherboard and DIMM traces
 - Calculates the round-trip time between MC and DIMM
 - Calculates discrepancies between different traces





- 4. Memory Controller sets DIMM's Mode Registers
 - Enable/disable features, fine tune timings







- CPU requests reads & writes
- Memory Controller schedules and optimizes memory accesses
- Memory Controller schedules periodic memory refreshes





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DRAM Functionality: Sleep (S3)



- CPU is powered off
- Memory controller is powered off
 - Clock Enable (CKE) memory signal is 0
 - The rest of the signals are in tri-state
- DIMM is in self-refresh state



DRAM Functionality: Waking up to SO



- CPU and Memory Controller are powered on
- Memory Controller pulls CKE to 1
- Read & Write requests can now be issued



Protocol Limitation



- No state information is held on the DIMM
 - e.g. calibration results, clock frequency
- No authentication between MC and DIMM
- Nothing prevents an attacker from impersonating a Memory Controller



Could we exploit that?

- Create a device that:
 - Attaches to the exposed signal pins
 - Impersonate a MC
- Extremely difficult when system is in S0
- Possible when system is in S3 sleep





Could we exploit that?

- Example:
 - The victim goes to a conference and leaves their system in sleep
 - The attacker attaches their rogue MC on the exposed pins
 - The attacker now owns the system





Requirements for Successful Exploitation

- Attach rogue MC to the victim system while in S3 sleep
 - Pull CKE to 1 to "wake up" the DIMM
 - Calibrate our rogue MC to the memory bus traces
 - Send memory read/write commands to the DIMM
 - Put CKE to 0 to put the DIMM to sleep
- Detach from the victim system and wake up normally



Signal Injection: CKE signal

- CKE must be pulled up above ~0.7V (logic high threshold)
- Find R_A so that DIMM sees logic high transition
- Larger values reduce current into victim MC (I_A)





Signal Injection: CKE signal





Rogue MC Calibration

- Need to train our rogue MC Victim with the new traces System
- Normally happens on first boot, can be done <u>any time</u>
- Calibration results will be the same on the same system + DIMM models





Adjust Clock Frequency

- Rogue MC provides a differential clock signal
- Attacker's clock can be different than the victim clock
 - Lower clock frequency \rightarrow Better signal integrity
- DDR4 has a default minimum frequency of 800 MHz
 - When "DLL-off" Mode Register is enabled, clock frequency can be significantly lower (as low as ~128 MHz)
 - This Mode Register is accessible to the attacker!



Issue Reads/Writes

- Follows JEDEC specification using off the shelf FPGA MC
- Iterate through reading all memory addresses
- Selectively write to memory addresses of interest





Rogue Memory Controller Board





Mechanical Connector





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Threat Model

- Requires a skilled HW attacker
- Physical access to the victim system
 - Invasive
 - Requires time to conduct
- May require prior knowledge of the victim system spec



Challenges / Future Work

- Complete mechanical connector
- Run attack against full length bus
- We anticipate signal noise/reflection challenges



Mitigations

- Short-term
 - Use hibernate (S4) instead of sleep (S3)
 - Focus on physical security
- Long-term
 - Enhance DIMM socket and motherboard design to restrict access to pins
 - Enhance JEDEC spec to perform authentication



Conclusions

- Described early research into a <u>physical attack</u> on DRAM memory
- A skilled attacker could exploit HW design and protocol limitations
- New line of research for physical security
- Short- and long-term mitigations are possible



Questions?

- Thank you!
- We are hiring!



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