

# Intercepting iCloud Keychain

Alex Radocea Longterm Security, Inc.



#### Alex Radocea

Co-founder at Longterm Security, Inc.

**Recovering CTF Addict** 

contact@longterm.io @defendtheworld



#### What is iCloud Keychain?



#### Secret Syncing & Recovery in the Cloud





#### Designed to be Highly Secure

- Strong end-to-end cryptography
- Resilient against a compromised backend, rogue insiders
- Resilient when an attacker has obtained a target's Apple ID password
  - Need an additional password or a trusted device



#### Critical Flaws Now Fixed

- We found critical flaws in undocumented, open-source compoments of the protocol
- Agenda: We'll describe previous work, how iCloud Keychain Syncing works, and the flaws in detail



#### Prior Work & Presentations Covering iCloud Keychain

- Andrey Belenko/ViaForensics <u>https://speakerdeck.com/belenko/on-the-security-of-the-icloud-keychain</u>
- Andrey Belenko/ViaForensics CVE-2015-1065 buffer overflows in keychain sync with MITM capability
- Ivan Krstic/é Behind the Scenes with iOS Security: Secret Synchronization - <u>https://www.blackhat.com/docs/us-16/materials/us-16-Krstic.pdf</u>
- Vladamir Katalov/Elcomsoft <a href="https://conference.hitb.org/">https://conference.hitb.org/</a> <a href="https://conference.hitb.org/">https://conference.hitb.org/</a>
- iOS 10 Security Guide https://www.apple.com/business/docs/ iOS\_Security\_Guide.pdf





#### **Robert Morris**

(Harry Naltchayan/THE WASHINGTON POST)

#### "Never underestimate the attention, risk, money and time that an opponent will put into reading traffic."



#### iCloud Keychain Components







HSM-Based Escrow System

















#### HSM-Based Escrow System Verification Passcode required SMS Verification





Sync'd Secrets Across All Trusted Devices

















# Sync'd Secrets<br/>Across All<br/>Trusted<br/>DevicesEnd to End<br/>EncryptionCircle of TrustApproval or<br/>Two-Factor/<br/>iCSC required<br/>to join



# iCloud Keychain Sync







#### Join Circle with Join Circle with Signed Syncing Circle **Apple ID** 256-bit ECDSA Apple ID **Establishes Password and Password and** on secp256r1 with SHA256 Trusted **Trusted Device** iCSC/Device Devices Approval Passcode



1. A creates a Circle





2. B requests to join

**Bob Application Ticket** 

Identity B Pubkey Identity B Sig

AppleID Password Signature



3. A approves





4. B countersigns





5. Third device approved by **B**...





6. Countersigned by all parties





# What happens when devices are lost while traveling?





#### iCloud Keychain Passwords Overview





#### How Does A New Device Join Without Approval?

- Circle does not reset when this happens
- Joining the circle requires a trusted device to sign the updated circle with an identity key...
- And Identity Keys not in the escrow
   kSecAttrAccessibleWhenPasscodeSetThisDeviceOnly

The class kSecAttrAccessibleWhenPasscodeSetThisDeviceOnly behaves the same as kSecAttrAccessibleWhenUnlocked, however it is only available when the device is configured with a passcode. This class exists only in the system keybag; they don't sync to iCloud Keychain, aren't backed up, and aren't included in escrow keybags. If the passcode is removed or reset, the items are rendered useless by discarding the class keys.



## Uncovering a hidden peer

- Undocumented, speculating this is for streamlining usability
- When a Circle is first established an "iCloud Identity" Key is also created as a "hidden" peer
  - Key is created with kSecAttrAccessibleWhenUnlocked, kSecAttrSynchronizable
- Available from iCloud Keychain Recovery
- Can be used to update the Syncing Circle, and trigger automatic coutersigning from all peers



#### Updated Circle Illustration -One Peer





#### Updated Circle Illustration -Two Peers





#### Which Backups Contain the Cloud Identity Key?

- Cloud Peer Backup sounds tricky, seems okay
- If available in iCloud Backup Keybags...
  - UID Key wrapping prevents Apple/Malicious Insider from accessing the data
- iCloud Keychain Escrow contains Cloud Identity Keys (kSecAttrSynchronizable)
  - Not available without SMS and either iCSC or passcode with two-factor authentication





#### End to End Encryption

#### Forward Secrecy & Deniability

ECDH Key Exchange, Verified with Peer Identity Keys

128-AES-CTR Encryption w/ Rotating Keys



#### iCloud Keychain Sync Transmits Data Across Apple Services





# E2E: Plaintext material only available on trusted devices

recent modification date will be synced. Items are skipped if the other member has the item and the modification dates are identical. Each item that is synced is encrypted specifically for the device it is being sent to. It can't be decrypted by other devices or Apple. Additionally, the encrypted item is ephemeral in iCloud; it's overwritten with each new item that's synced.



## **OTR KEX Messages**

#### Initiator

#### Receiver



# **OTR KEX Messages**

Receiver

Peer Identity Keys from SOSCircle used for Signature Verification of Ephemeral DH Keys

Initiator



# **OTR KEX Messages**

Initiator

Receiver

Secure Channel used to establish long-term keys, exchange messages, and ultimately passwords. No further encryption of passwords at this point



#### Pairwise, Fanout Negotiation







#### Apple's iCloud Keychain Security Goals

- "Sync passwords between iOS devices and Mac computers without exposing that information to Apple"
- Also protect password material:
  - When the iCloud account is compromised
  - When iCloud is compromised by a rogue insider or external attackers
  - When third parties access user accounts







#### iCloud Keychain Sync Remote Attack Graph





#### **OTR Flaws**

- CVE-2017-2448 OTR Cryptographic Failure
- CVE-2017-2451 OTR Memory Corruption
- Exacerbated by lack of TLS key pinning on KVS communications



#### CVE-2017-2448 - Goto Fail Redux



```
result = ReadLong(signatureAndMacBytes, signatureAndMacSize,
&xbSize); [1]
```

```
require_noerr(result, exit);
require_action(xbSize > 4, exit, result = errSecDecode);
```

```
require_action(xbSize <= *signatureAndMacSize, exit, result =
errSecDecode);</pre>
```

```
uint8_t signatureMac[CCSHA256_OUTPUT_SIZE];
```

```
cchmac(ccsha256_di(), sizeof(m2), m2, xbSize + 4,
encSigDataBlobStart, signatureMac);
```

require(xbSize + kSHA256HMAC160Bytes <= \*signatureAndMacSize, exit);
[2]</pre>



...

result = ReadLong(signatureAndMacBytes, signatureAndMacSize, &xbSize) [1]

```
require_noerr(result, exit);
require_action(xbSize > 4, exit, result = errSecDecode);
```

```
require_action(xbSize <= *signatureAndMacSize, exit, result =
errSecDecode);</pre>
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uint8\_t signatureMac[CCSHA256\_0UTPUT\_SIZE];

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```

require(xbSize + kSHA256HMAC160Bytes <= \*signatureAndMacSize, exit);
[2]</pre>



```
result = ReadLong(signatureAndMacBytes, signatureAndMacSize,
&xbSize): [1]
require_noerr(result, exit);
require_action(xbSize > 4, exit, result = errSecDecode)
require_action(xbSize <= *signatureAndMacSize, exit, result =</pre>
errSecDecode);
uint8_t signatureMac[CCSHA256_OUTPUT_SIZE];
cchmac(ccsha256_di(), sizeof(m2), m2, xbSize + 4,
encSigDataBlobStart, signatureMac);
require(xbSize + kSHA256HMAC160Bytes <= *signatureAndMacSize, exit);</pre>
[2]
```



result = ReadLong(signatureAndMacBytes, signatureAndMacSize, &xbSize) · [1]

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uint8\_t signatureMac[CCSHA256\_0UTPUT\_SIZE];

```
cchmac(ccsha256_di(), sizeof(m2), m2, xbSize + 4,
encSigDataBlobStart, signatureMac);
```

require(xbSize + kSHA256HMAC160Bytes <= \*signatureAndMacSize, exit);
[2]</pre>

#### CVE-2017-2448 - Goto Fail Redux

static OSStatus SecVerifySignatureAndMac(SecOTRSessionRef session, bool usePrimes, const uint8\_t \*\*signatureAndMacBytes, size\_t \*signatureAndMacSize) {

```
OSStatus result = errSecDecode;
```

result = ReadLong(signatureAndMacBytes, signatureAndMacSize, &xbSize); [1]

require\_noerr(result, exit);
require\_action(xbSize > 4, exit, result = errSecDecode);

require\_action(xbSize <= \*signatureAndMacSize, exit, result =
errSecDecode);</pre>

uint8\_t signatureMac[CCSHA256\_OUTPUT\_SIZE];

cchmac(ccsha256\_di(), sizeof(m2), m2, xbSize + 4, encSigDataBlobStart, signatureMac);

require(xbSize + kSHA256HMAC160Bytes <= \*signatureAn MacSize, exit);
[2]</pre>

...

exit:

bzero(m1, sizeof(m1)); bzero(m2, sizeof(m2)); bzero(c, sizeof(c));

return result;

- Error handling erroneously returns successfully on parsing failure
- Encoding an invalid size in an OTR packet establishes a DH key exchange and bypasses signature verification



#### CVE-2017-2448 - Sample Trigger in 32 Bytes

00 0x00 0x02 0x12 0x00 04 0x00 0x00 0x00 0x18 08 0x41 0x41 0x41 0x41 0C 0x1c 0x41 0x41 0x41 0x41

int i = 0; payload[i++] = 0x00; payload[i++] = 0x02; //version 2

payload[i++] = kSignatureMessage; // packet type

payload[i++] = 0; //xbsize
payload[i++] = 0; //xbsize
payload[i++] = 0; //xbsize
payload[i++] = N-8; //xbsize

payload\_length = N;

?OTR:AAISAAAAGEFBQUFBQUFBQUFBQUFBQUFBQUFBQUFBQUE=.



#### Signature Bypass Attack Impact

- MITM Attacker could impersonate existing peers to negotiate secrets
- OTR protocol encrypts using ephemeral keys, verified with the peer identity keys
- Silent attack on targets with 100% reliability



#### Apple's iCloud Keychain Security Goals (without OTR fix)

- "Sync passwords between iOS devices and Mac computers without exposing that information to Apple"
- Also protect password material:
  - When the iCloud account is compromised
  - When iCloud is compromised by a rogue inside external attackers
  - When third parties access user accounts



#### CVE-2017-2451 - Stack Clash

```
result = ReadLong(signatureAndMacBytes,
signatureAndMacSize, &xbSize);
```

uint8\_t xb[xbSize];

- Same Routine as CVE-2017-2448
- MITM attacker controls stack allocation size
- Long OTR packet results in data being allocated in adjacent thread's stack



#### Stack Overlap Attack Impact

- Potential sandbox escape into secd (as root)
- Malicious local application could potentially gain access to device keychains
- Remotely triggerable as well
- Tricky to exploit due to guard pages, trigger races against a crash



# Wrapping up

- Exciting to see strong and usable end-to-end encryption for the masses
- We covered the Keychain Sync Protocol in depth
- We reviewed a critical vulnerability in OTR that undermined the End to End Encryption



#### Next Steps for the Security Industry

- Should this have been discovered after Goto Fail?
  - Strikingly similar, same code project.
  - See Crypto Testing Talk
- Are the protocol details sufficiently transparent to users?
  - Mostly open source, but we're still the first to discuss OTR publicly
- More research needed on the two-factor implementation, and its interface with iCloud Keychain Recovery and iCloud Keychain Syncing



#### **Questions?**



## Appendix



#### **Circle Protocol Parameters**

- Apple ID Password converted to ECC keypair using PBKDF2 and X9.63
- Identity Keys are 256-bit keys on the secp256r1 curve
  - Stored in Keychain with kSecAttrAccessibleWhenUnlockedThisDeviceOnly protection class
  - Cloud Identity Key
     kSecAttrAccessibleWhenUnlocked and synchronizable



#### **OTR Encryption Parameters**

- NIST Curve (secp256r1)
- ECDH with ephemeral keys over secp256r1
- ECDSA signatures over secp256r1 with SHA-256
- SHA256-HMAC-160
- 128-bit AES-CTR used for encryption



#### OTR Asynchronous Key Exchange

BOB ALIC	E
$\operatorname{AES}_r(g^x), \operatorname{HASH}(g^x)$ 1 "D-H Commit Message"	<ol> <li>Hash commitment</li> <li>Diffie-Hellman Key Exchange</li> </ol>
g <sup>y</sup> 2 "D-H Key Message"	<ul> <li>Build Freihaus (Key Exercise)</li> <li>3. Encrypted exchange of long-term keys &amp; signatures</li> </ul>
$\begin{split} M_B &= \mathrm{MAC}_{K_{m_1}}(g^x, g^y, pub_B, keyi \\ X_B &= \{pub_B, sig_B(M_B)\} \end{split}$	$d_B)$
$r$ $AES_c(X_B), MAC_{K_{m_2}}(AES_c(X_B))$ "Reveal Signature Message"	
$\begin{split} M_{A} &= \operatorname{MAC}_{K_{m'_{1}}}(g^{y}, g^{x}, pub_{A}, keyie_{A'}, X_{A} &= \{pub_{A'}, keyid_{A'}, sig_{A}(M_{A})\} \end{split}$	$d_A$ (3)
$\operatorname{AES}_{c'}(X_A), \operatorname{MAC}_{K_{m'_2}}(\operatorname{AES}_{c'}(X_A))$ "Signature Message"	)

https://blog.cryptographyengineering.com/2016/03/21/attack-of-week-apple-imessage/

