Timing Attacks Have Never Been So Practical: Advanced Cross-Site Search Attacks





About me: Nethanel Gelernter

- Security Researcher / Hacker
 - Web application security
 - Ph.D., hacks, research papers, talks, etc.
- Cyberpion



- Exploring new attack vectors & developing defenses against them
- Leading the cyber-security studies & research in the College of Management, Israel



The College of Management Academic Studies

Agenda – practical timing attacks

- Cross-site search (XS-search) attacks & Response inflation
- Challenges

– When response inflation is impossible

- Browser-based XS-search attacks
- Second-order XS-search attacks



Cross-Site Search Attacks

- Gelernter & Herzberg, CCS' 2015
- Exploit 'search' timing side-channel
- 'Search' in private-data kept by web-service
- Practical:
 - Tested on hundreds of Gmail users
 - Real world conditions
- Example: find user name
 - From lists of 2000 common (first and last) names

CYBERPA

– Takes about a minute

Cross-site attacker model

- Main model for web attacks
- The victim's browser is authenticated to services that hold private records (e.g., Gmail)
- The victim visits the attacker's website





Cross-site attacker model

- Cross-site search over user's data in service
 - Attacker cannot access the content of the response
 - Same Origin Policy
 - The attacker can measure the response *time* (*T*)



CYBERPA

XS-Search example: user name

- Find out whether the user is Alice or Bob...
- Compare:
 - T(Bob): response time for 'messages sent by Bob'
 - T(Alice): response time for 'messages sent by Alice'



What else can XS-Search expose?

Structured information





Contacts



Nethanel Gelernter



Email content

Outlook

15



YAHOO!

XS-Search: Basic Flow

• Find the answer for a Boolean question

- Three steps:
 - Transform the question into a search request
 - Send search requests and collect samples
 - Analyze response times \rightarrow answer the question!



XS-Search: Basic Flow – 1st Step

- Is the name of the user *Alice*?
 - in:sent from:Alice
- Is she related to bob@gmail.com?
 - bob@gmail.com&st=100
- Does Alice have an affair with Charlie
 - "I love you" to:Charlie from:Alice



XS-Search: Basic Flow – 2nd Step

- Send a Challenge request
 - Is the user name Alice?
 - True: a Full response is returned (has some content)
 - False: an empty response is returned





XS-Search: Basic Flow – 2nd Step

- Send a Dummy request
 - Is the user name *fdjakdhasd*?
 - The response is expected to be **empty**



CYBERPA

XS-Search: Basic Flow – 2nd Step



CYBERPON

XS-Search: Basic Flow – 3rd Step





Practical timing attacks: challenges

- Timing attacks
 - Delays depend on dynamically-changing factors, e.g.:
 Congestion and concurrent processes in client and server
- Practical attacks
 - Minimal time
 - Exploit also short visits of users
 - Minimal number of requests
 - Avoid detection and blocking
 - E.g., by server's anti-DoS defenses



Response Inflation

- Increase the size difference between full and empty responses
- Larger difference in size → Larger difference in time



Larger \rightarrow Slower

CYBERP^AON

Response Inflation

- Search requests have many parameters
- Some of them are reflected in the responses as a function of the number of results

https://example.com/search?reflected_parameter=value



Empty response

value	value	value
value	value	value
value	value	value

Full response



Response Inflation

 Sometimes, the attacker send very long strings as the value of the reflected parameter

https://example.com/search?reflected_parameter=Long string

Long string.....

Empty response

Long string.....Long string.....

Full response



Response inflation example

- Exploiting Gmail search in the HTML view
- The query itself!
 - Appears once for each entry (50 max by default)
 - Can be inflated to 8KB
- Up to 400KB response size inflation!



Terms - Privacy - Gmail Blog - Google Home



But...





What if there is no response inflation?





What if there is no response inflation?

- Browser-based XS-search
 - When there is **some** difference in the response size
- Second-order XS-search
 - When there is **no** difference in the response size!



Browser-based (BB) XS-Search

- Statistical tests and divide and conquer algorithms
 - Gelernter & Herzberg, CCS' 2015
- Browser-based timing side channel
 - Van Goethem et al.,CCS' 2015
- Algorithmic improvements



Classical vs. BB timing attacks

- Classical timing attacks:
 - Load the resources from the server several times to collect time measurements

- Browser-based timing attacks:
 - Load all the resources from the server once and cache them
 - Then load them from the cache many times to collect time measurements



Classical vs. BB timing attacks

- Exploiting / measurements affected by
 - Classical: network delay, server processing time, browser processing time
 - Browser-based: browser processing time

- Can be used to differentiate between
 - Classical: large/small resources, high/low server processing time
 - Browser-based: large/small resources



BB XS-Search: Basic Flow

• Find the answer for a Boolean question

 Changing only the second step of the original XS-Search attack



BB XS-Search: Basic Flow – 2nd Step

- Send a Challenge request
 - Is the user name Alice?
 - True: a Full response is returned (has some content)
 - False: an empty response is returned





BB XS-Search: Basic Flow – 2nd Step

- Send a **Dummy** request
 - Is the user name *fdjakdhasd*?
 - The response is expected to be empty





Nethanel Gelernter

CYBERPAON

Browser-based (BB) XS-Search

- Algorithmic improvements
- Not for Boolean questions
 - Basic flow only Boolean questions
 - Is the victim's name Alice?
- Answering multiple choice questions
 - E.g., which names out of many options are matching the victim?
- Optimally use the browser-based timing sidechannel



Browser-based (BB) XS-Search

- Evaluation compared to both the previous works
- Repeating attacks/experiments done in each of them
 - Original XS-Search: extract victim's names from Gmail
 - BB timing attacks: extract victim's age from Facebook
- Significant improvement!
- In this talk: only one example



- Gmail example
 - The goal of the attacker: extract the first and last names of the victim out of a list of 2000 names
 - XS-Search results:
 - 90% success rate (both first and last name found)
 - 1 minute on average
 - 2.6% false positive



- How to answer multiple-answer questions efficiently?
- The optimized multiple term identification (OMTI) algorithm
 - Divide and conquer algorithm
 - Relying on the OR operator
 - Different dummy search request is sent every round



- Rely on browser-based timing side-channel to optimize the OMTI algorithm
- Observation: empty responses are (almost) identical
 - No need to send dummy requests in every round
 - No need to reload the empty response in every round
 - Rely on previous measurements!



- Evaluation of the attack on 5 different Gmail accounts
 - 15-16 times on each of them
- Significant improvement!
 - 41.6 seconds on average (compared to 1 minute)
 - 92.3% success (compared to 89.7%)
 - 1.3% false positive (compared to 2.6%)



• DEMO



• The problem: sometimes the size difference is negligible

• For example: a sentence that appears in a single email



value

Full response



- Second-order attacks
 - First, manipulate the attacked web application
 - Make it (more) vulnerable
 - Exploit the vulnerability
- Second-order XS-search attacks
 - First manipulate the attacked storage
 - Create significant response inflation
 - Launch browser-based XS-search attack

CYBERP

- Two SO XS-search attacks
 - -Simple
 - –Inflating



- Model
 - Storage
 - Many records
 - A secret appears in one of the records

- Attacker can manipulate the storage remotely
 - E.g., email accounts
 - Another example later...

CYBERP^AON

Simple SO XS-Search attack

• The problem: the secret appears only once in the storage

• Simple solution: the attacker will add additional records that contain the secret!



Simple SO XS-Search attack



CYBERPON

Simple SO XS-Search attack

• Example: extracting Facebook password-reset code from Yahoo! email



- Creates significant response inflation effect
 - Increase the size difference between empty and full response

 Unlike all the previous attacks: the empty response will be (significantly) larger than the full response



- The challenge of the attacker:
 - Find a secret out of a large dictionary of possible values
- Notations
 - *M* maximal number of results
 - Match-all record a record that contains all the possible values for the secret
 - *Inflating record* a record that significantly inflates the size of every response containing it



• Attack process

<u>First part:</u>

- Plant one *match-all inflating* record in the storage
- Plant additional *M*-1 *match-all* records
- Additional record(s) may be added as a result of the victim's operations, or via operations triggered by the attacker

Second part:

– Launch BB XS-search attack!





CYBERPAON

- Inflating record in email service providers
 - Email headers
 - From
 - То



- Example: extracting Visa/Mastercard credit card number
 - Structured information
 - VVVV-XXXX-YYYY-ZZZZ

- First and last names: extract 2 out of 2000
 Done successfully!
- Credit card number: extract 4 out of 10000
 - Should not be much harder



- Example: extracting Visa/Mastercard credit card number
- Match-all record a record that contains all the possible 4-digit sequences

– Possibly as an attachment

 Inflating match-all record – a match-all record with very long From field



- Gmail example
- How?
 - Cross-site search requests are now blocked in both the HTML and standard views
- Cross-site search attack without sending cross-site search requests?



- Gmail example
- Exploiting the autocomplete feature!



• Gmail example: the manipulated storage





• Gmail example: full response (size is small)



CYBERPAON

• Gmail example: empty response (size is very large)





• DEMO



- Evaluation results
 - 96% success rate within less than 50 seconds
 - Yet, in the other 4% percent, 3 out of 4 sequences were found, and it was possible to detect the error and to fix it



• The challenge: manipulations on the storage can be detected!

• Solution: manipulate the storage in a way that will not be detected by the user

• HOW?



- Emails solution: abuse anti-spam mechanisms
- The planted emails will be marked as spam

 Users do not get notifications for spam emails
 - Users (usually) do not visit their spam folder
- Only when it is possible to search in the spam and in the other folders using the same request
 - E.g., Gmail
 - in:inbox OR in:spam



- Search history
- Two requirement for inflating SO XS-Search attack:
 - -Inject records to the search history log
 - DONE: Gelernter & Grinstein & Herzberg, ACSAC 2015
 - Inject an inflating record



Search History

Bing • example: inflating SO **XS-Search** attack to extract search history

search C	All types \checkmark All dates $\overline{II} \checkmark$ Show all	Off On Clear all
TODAY		
search0		9:52 am 🛛 🗙
search1	Areas of Specialty SearchOne search1.com/areas-specialty Homepage SearchOne search1.com Medicare Coverage Database – Centers cms.gov/medicare-coverage-d	9:52 am 🛛 🗙
search2	Cook County Sheriff Inmate Locator www2.cookcountysheriff.org/	9:52 am 🛛 🗙
search3	The Hoppin 5k - RaceIt.com www.raceit.com/search3/even Compass Micro - An Authorized Epson a www.compassmicro.com/parts search3.com (@search3) Twitter twitter.com/search3	9:51 am 🛛 🗙
search4		9:51 am ×
search5	Search Results - Magellan GPS www.magellangps.com/Search FVP - Search Results From: search5.fypimageviewer.com	9:51 am 🛛 🗙



Defenses (briefly)

• If possible - blocking cross-site search requests

- In other cases make it harder to exploit
 - Block inflation techniques
 - Rate limit
- Like (almost) every other web-application attack the challenge is to find all the vulnerable spots



Conclusions

- Advanced cross-site search attacks
 - Browser-based
 - Second order
- Practical!
- Many vulnerable websites
 - Including popular ones



Thank you!

CYBERPAON

Questions?

