ROPInjector: Using Return-Oriented Programming for Polymorphism and AV Evasion

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• Return Oriented Programming (ROP) has been used in the past explicitly for **DEP evasion** in software exploitation scenarios.

• We propose a completely **different use of ROP**!

• We propose ROP as a **polymorphic alternative** to achieve **AV evasion**!

\[ \text{1 PE} + \text{1 shellcode} \rightarrow \text{M different variations} \]
Local infection of benign PE executables with well-known alarming malicious code (i.e., shellcode)
Why ROP for AV evasion?

• **Does not raise suspicious** ➔ Borrowed code (i.e., rop gadgets) is of course benign. The only possible detection footprint is the instructions (i.e., typically push) that insert the addresses of the ROP gadgets into the stack.

• **Generic** ➔ ROP can be used to transform any given malware shellcode to a ROP-based equivalent.

• **Polymorphism** ➔ Use different ROP gadgets or use same ROP gadgets found in different address
1. The new resulting PE should **evade antivirus detection**

2. The benign PE should not be **corrupted/damaged**

3. The tool should be generic and completely automated

4. Should not require a writeable code section to mutate (i.e., execute ROP chain)

**Challenges Accepted!**
A quick historical overview

plain malware code
\x59\xE8\xFF\x6B\x5F\xFF\x6A\x0F\x59\xE8\xFF

string signatures
\x6B\x5F\xFF\x6A\x0F
A quick historical overview

plain malware code

simple obfuscation (NOPs/dead-code in-between)


\x59\xE8\xFF\x6B\x5F\x90\xFF\x90\x6A\x0F\x59\xE8

\x5B\x90*\xFF*\x6A\x0F

string signatures

regex signatures

variability
A quick historical overview

- Plain malware code
- Simple obfuscation (NOPs/dead-code in-between)
- Oligomorphism

Decoder encoded payload

PC
\x6A\x0F\x59\xE8\x00\xFF \x6B\x5F**************

string signatures
regex signatures
static analysis (disassembly, CFGs)

if RWX and performs then alarm
A quick historical overview

- Plain malware code
- Simple obfuscation (NOPs/dead-code in-between)
- Oligomorphism
- Self-modifying code
- Metamorphism

- String signatures
- Regex signatures
- Static analysis (disassembly, CFGs)
- Dynamic analysis (emulation, sandboxing, behavior-based signatures)

```
push eax
mov [esp-4], eax
sub esp, 4
```
1. Shellcode analysis

2. Find ROP gadgets in PE

3. Transform shellcode to an equivalent ROP chain

4. Inject in PE missing ROP gadgets (if required)

5. Assemble ROP chain building code in PE

6. Patch the chain building code into the PE
• The analysis of the shellcode aims to obtain various information so that it can safely replace shellcode instructions with gadgets

• In particular, for each instruction, ROPInjector likes to know:
  – what registers it reads, writes or sets
  – what registers are free to modify
  – its bitness (a mov al, X or a mov eax, X ?)
  – whether it is a branch (jmp, conditional, ret, call)
  – and if so, where it lands
  – whether it is a privileged instruction (e.g. sysenter, iret)
  – whether it contains a VA reference
  – whether it uses indirect addressing mode (e.g. mov [edi+4], esi)
• Scaled Index Byte (SIB) enables complex indirect addressing modes
  – e.g. \texttt{mov [eax+8*edi+10], ecx}

• We want to avoid SIBs in the shellcode because they are:
  – long: >3 bytes \(\rightarrow\) unlikely to be found in gadgets
  – rarely reusable
  – reserve at least 2 registers
• ROPInjector performs a technique that we call it **Unrolling of SIBs** to transform them into simpler instructions

```assembly
mov eax, [ebx+ecx*2]  # Unrolled
sal eax, 1
add eax, ebx
mov eax, [eax]
```

• **ecx is freed at this point**
• **shorter instructions**
• **reusable gadgets** (either found or injected)

• With this technique we achieve
  – increased chances of finding suitable gadgets
  – less gadgets being injected
STEP 2: Find ROP Gadgets in PE (1/2)

1. First find returns of type:
   - ret(n) or
   - pop regX
     jmp regX or
   - jmp regX
   • used only if a “loader” gadget of the following form is also found:
     pop regX
     [any of the first 2 endings above]

2. Then search backwards for more candidate gadgets
ROPInjector automatically resolves redundant instructions in ROP gadgets, in order to avoid errors during execution of ROP code.

- Maximize reusability of ROP gadgets

- Avoid injecting unsafe ROP gadgets
  - modify non-free registers
  - are branches
  - write to the stack or modify esp
  - are privileged
  - use indirect addressing mode
STEP 3: Transform shellcode to ROP chain

• First, we translate shellcode instructions to an Intermediate Representation (IR).
• Next we translate ROP gadgets found in PE to an IR.
• Finally, a mapping is performed between the two IRs
  – 1 to 1
  – 1 to many
## STEP 3: Intermediate Representation

<table>
<thead>
<tr>
<th>IR Type (20 in total)</th>
<th>Semantics</th>
<th>Eligible instructions</th>
</tr>
</thead>
</table>
| ADD_IMM              | regA += imm    | add r8/16/32, imm8/16/32  
|                      |                | add (e)ax/al, imm8/16/32  
|                      |                | xor r8/16/32, 0  
|                      |                | cmp r8/16/32, 0  
|                      |                | inc r8/16/32  
|                      |                | test ra32, rb32 (with ra == rb)  
|                      |                | test r8/16/32, 0xFF/FFFF/FFFFFFFF  
|                      |                | test (e)ax/al, 0xFF/FFFF/FFFFFFFF  
|                      |                | or ra32, rb32 (with ra == rb)  |
| MOV_REG_IMM          | mov regA, imm  | mov r8/16/32, imm8/16/32  
|                      |                | imul r16/32, r16/32, 0  
|                      |                | xor ra8/16/32, ra8/16/32  
|                      |                | and r8/16/32, 0  
|                      |                | and (e)ax/al, 0  
|                      |                | or r8/16/32, 0xFF/FFFF/FFFFFFFF  
|                      |                | or (e)ax/al, 0xFF/FFFF/FFFFFFFF  |
STEP 3: Mapping examples

• 1-1 mapping example
  – Shellcode:
    mov eax, 0  → MOV_REG_IMM(eax, 0)
  – Gadget in PE:
    and eax, 0  → MOV_REG_IMM(eax, 0)
  → MOV_REG_IMM(eax, 0)

• 1-many mapping example
  – Shellcode:
    add eax, 2  → ADD_IMM(eax, 2)
  – Gadget in PE:
    inc eax  → ADD_IMM(eax, 1)
    ret
  → ADD_IMM(eax, 1)
STEP 4: Gadget Injection

- In some cases PE does not include the required ROP gadgets

- Simply injecting ROP gadgets would raise alarms using **statistics**
  - Presence of successive ret instructions

- To this end, ROPInjector inserts ROP gadgets **scattered**, and in a benign looking way to avoid alarms
  - \(0xCC\) caves in .text section of PEs (padding space left by the linker)
  - Often preceded by a RET (due to function epilogue)
• Assuming missing gadget is `mov ecx, eax` and we find the following:

    0xCC cave:

    <other instructions>

    epilogue:
    mov esp, ebp
    pop ebp
    return:
    ret(n)
    CCCCCCCCCCCCCCCCCCCCCCCCC
STEP 4: Gadget Injection

- Assuming missing gadget is `mov ecx, eax` and we find the following 0xCC cave:

```assembly
<other instructions>
jmp epilogue
mov ecx, eax
jmp return
epilogue:
    mov esp, ebp
    pop ebp
return:
    ret(n)
CCCCCCCCC
```
• Assuming missing gadget is `mov ecx, eax` and we find the following 0xCC cave:
• Assuming missing gadget is `mov ecx, eax` and we find the following 0xCC cave:

```assembly
<other instructions>
jmp epilogue
mov ecx, eax
jmp return
epilogue:
    mov esp, ebp
    pop ebp
return:
    ret(n)
```

CCCCCCCCC
STEP 5 and 6: Assemble and patch the ROP chain into PE

• In step 5 we insert the code that loads the ROP chain onto the stack (mainly PUSH instructions)

• In step 6 we patch the new PE → ROPInjector extends the .text section (instead of adding a new one that would raise alarm) and then goes on to repair all RVAs and relocations in the PE.

• ROPInjector includes two different methods to pass control to the ROPed shellcode
  – Run first
  – Run last
STEP 6: PE Patching (1/2)

**Before injection**
- NT Header
- Section Header (.text)
- Section Header (.data)
- Section Header (.rsrc)
- .text
- .data
- .rsrc

**After Injection**
- NT Header
- Section Header (.text)
- Section Header (.data)
- Section Header (.rsrc)
- .text
- .data
- .rsrc

- NT header checksum recalculated
STEP 6: PE Patching (2/2)

Run first:

1. NT Header
   AddressOfEntryPoint

2. Section .text
   jmp-to-malware

3. [malware code]
   [replaced code]
   jmp-back

Run last:

(very good anti-emulation results)

Previous calls to
ExitProcess() / exit()

(jmp-to-malware)

Section .text

[malware code]
ExitProcess()
Evaluation

- Implemented in native Win32 C
- 9 32bit executables
  - firefox.exe, java.exe, AcroRd32.exe, cmd.exe, notepad++.exe and more
- Various combinations
  - No patching at all
  - ROP’ed shellcode and run last
  - Intact shellcode passed control during exit (run last)
  - ROP’ed shellcode and delayed execution (20 secs via Win32 Sleep())
  - Intact shellcode
- 2 of the most popular payloads of MSF
  - reverse TCP shell
  - meterpreter reverse TCP
- VirusTotal
  - at the time it employed 57 AVs
Evasion rate: reverse TCP shell

Evasion ratio

<table>
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<tr>
<th>Original file</th>
<th>ROP-Exit</th>
<th>Exit</th>
<th>ROP-d20</th>
<th>Shellcode</th>
</tr>
</thead>
<tbody>
<tr>
<td>AcroRd32.exe</td>
<td>Acrobat.exe</td>
<td>cmd.exe</td>
<td>Rainmeter.exe</td>
<td>firefox.exe</td>
</tr>
</tbody>
</table>
Evasion rate: meterpreter reverse TCP

Evasion ratio

- Original file
- ROP-Exit
- Exit
- ROP-d20
- Shellcode

Experiments were conducted with various applications:
- AcroRd32.exe
- Acrobat.exe
- cmd.exe
- Rainmeter.exe
- firefox.exe
- java.exe
- wmplayer.exe
- nam.exe
- notepad++.exe

The evasion ratio was measured across these applications, showing varying degrees of success in evading detection.
Overall evasion results

- 100% most of the times
- 99.31% on average
• Current signature-based detection methods are no longer effective
  – we shown that by using ROP we can reduce the footprint to benign stack modifying instructions

• Behavioral analysis is tough to perform exhaustively
  – we shown how to easily bypass it by running right before process exit

• “Default distrust all” policy
  – Checksums and certificates is the poor user’s last line of defense at the moment