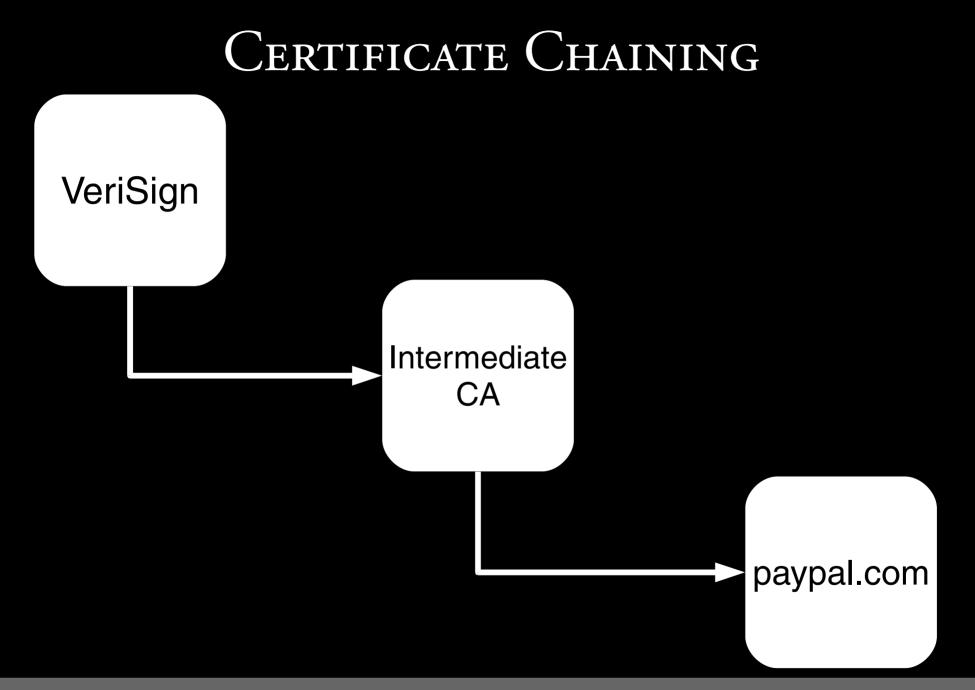
More Tricks For Defeating SSL In Practice

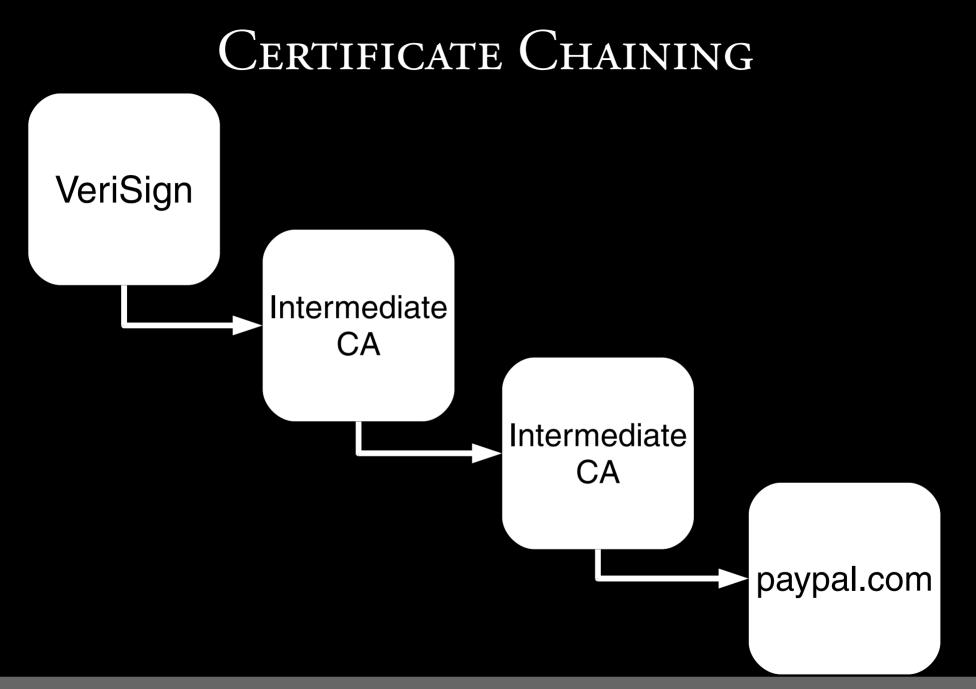
Moxie Marlinspike moxie@thoughtcrime.org



Once Again, The Back Story...

In the past, I've talked about BasicConstraints...



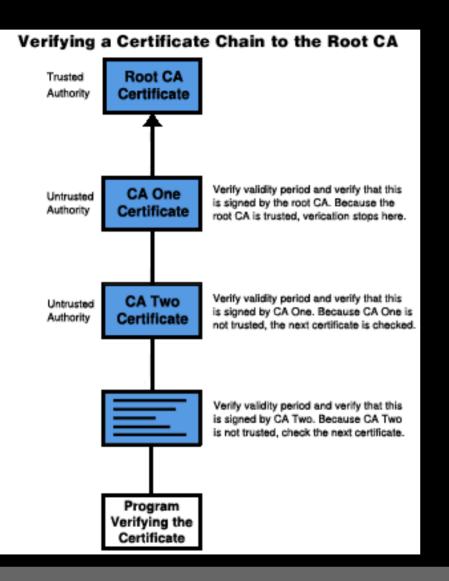


How do we verify these things?

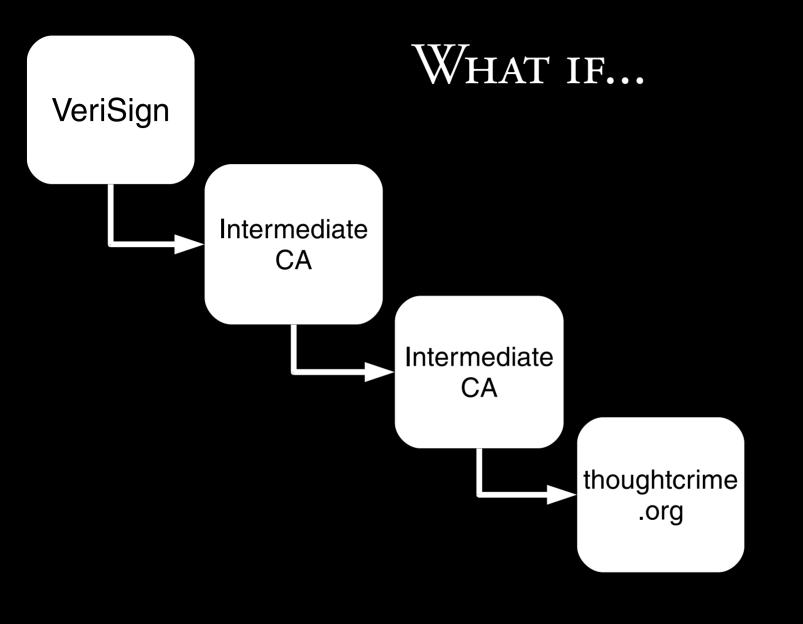


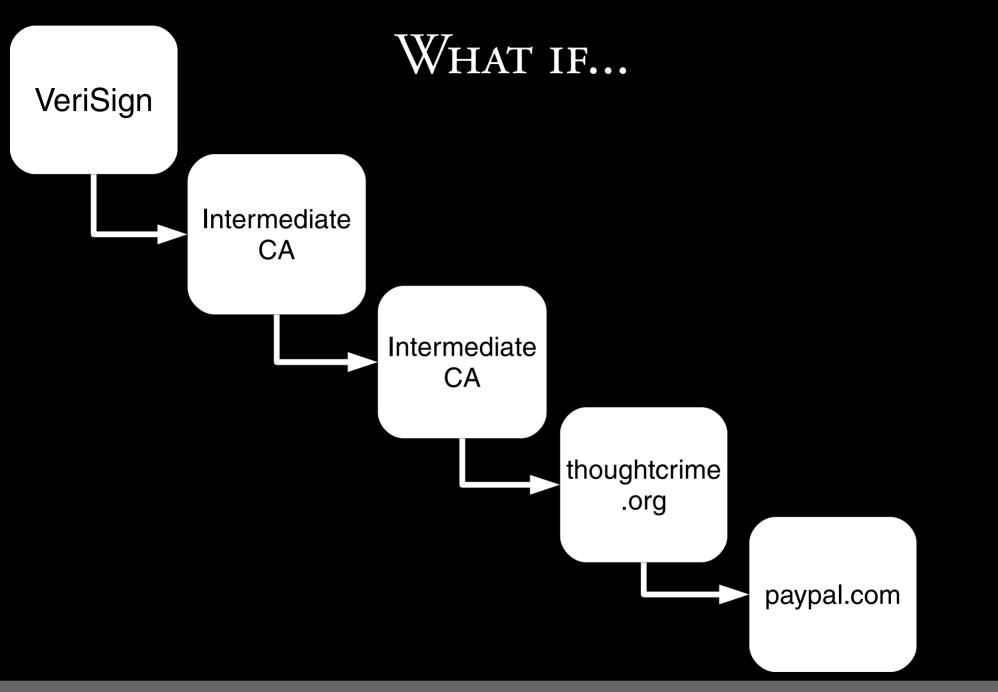
- Verify that the name of the leaf node is the same as the site you're connecting to.
- Verify that the leaf certificate has not expired.
- Check the signature.
- If the signing CA is in our list of trusted root CAs, stop. Otherwise, move one up the chain and repeat.

Here Be Dragons



- Very tempting to use a simple recursive function.
- Everyone focuses on the signature validation.
- The result of a naïve attempt at validation is a chain that is complete, but nothing more.







- Verify that the name of the leaf node is the same as the site you're connecting to.
- Verify that the leaf certificate has not expired.
- Check the signature.
- If the signing CA is in our list of trusted root CAs, stop. Otherwise, move one up the chain and repeat.

Something must be wrong, but...

- All the signatures are valid.
- Nothing has expired.
- The chain is in tact.
- The root CA is embedded in the browser and trusted.

BUT WE JUST CREATED A VALID CERTIFICATE FOR PAYPAL, AND WE'RE NOT PAYPAL?

THE MISSING PIECE...

...IS A SOMEWHAT OBSCURE FIELD.

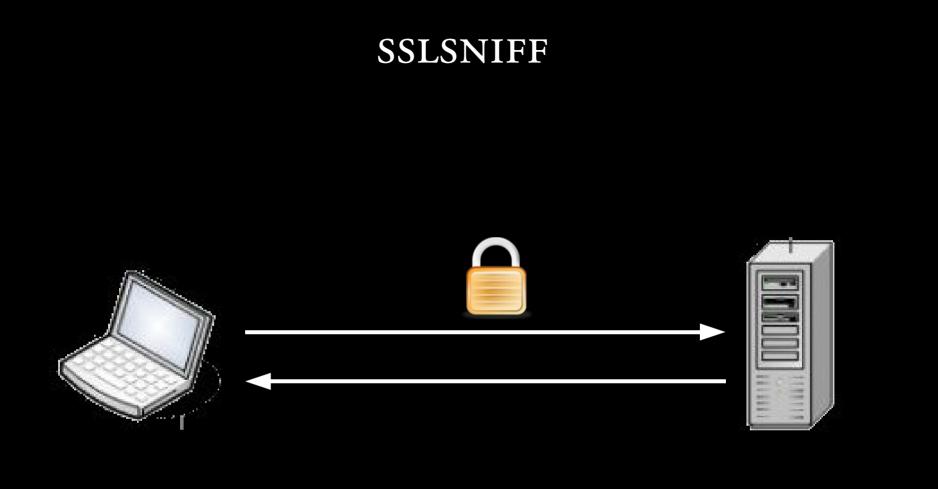
<u>F</u> ile <u>E</u> dit	<u>V</u> iew <u>T</u> erminal Ta <u>b</u> s <u>H</u> elp
moxie@sea	arching: ~/Desktop/b 💥 moxie@searching: ~/Desktop/b 💥 moxie@searching: ~/Desktop/b 💥
	f8:c9:0f:24:d2:c7:c2:92:0c:13:54:93:d5:9b:c7: 0e:fa:19:a8:d5:d3:f7:ab:5d Exponent: 65537 (0x10001) X509v3 extensions: X509v3 Key Usage: critical Digital Signature, Non Repudiation, Key Encipherment, Data Encip
herment	X509v3 Subject Key Identifier: DF:48:EF:25:BF:D2:23:B0:F0:C2:AC:FA:5A:85:50:74:FF:F9:34:EF X509v3 CRL Distribution Points: URI:http://crl.geotrust.com/crls/globalca1.crl
с	X509v3 Authority Key Identifier: keyid:BE:A8:A0:74:72:50:6B:44:B7:C9:23:D8:FB:A8:FF:B3:57:6B:68:6
	X509v3 Extended Key Usage: TLS Web Server Authentication, TLS Web Client Authentication X509v3 Basic Constraints: critical
	CA:FALSE ature Algorithm: sha1WithRSAEncryption 7a:58:f9:88:14:cb:77:32:aa:83:12:de:d9:15:74:8e:34:e3: 66:ca:bc:24:2c:28:96:54:cd:be:51:56:60:87:e3:be:c6:2e: 86:7e:74:c1:68:01:b6:8c:07:c6:a2:0c:a4:36:ca:e1:a8:e9:

BACK IN THE DAY

- Most CAs didn't explicitly set basicConstraints: CA=False
- Whether the field was there or not, most SSL implementations didn't bother to check it.
- *Anyone* with a valid leaf node certificate could create and sign a leaf node certificate for *any other* domain.
- When presented with a complete chain, IE, Outlook, Konqueror, OpenSSL, and others considered it valid...

AND THEN IN 2002...

- Microsoft did something particularly annoying, so I blew this up by publishing it.
- Microsoft claimed that it was impossible to exploit.
- So I also published the tool that exploits it.



SSLSNIFF



SSLSNIFF

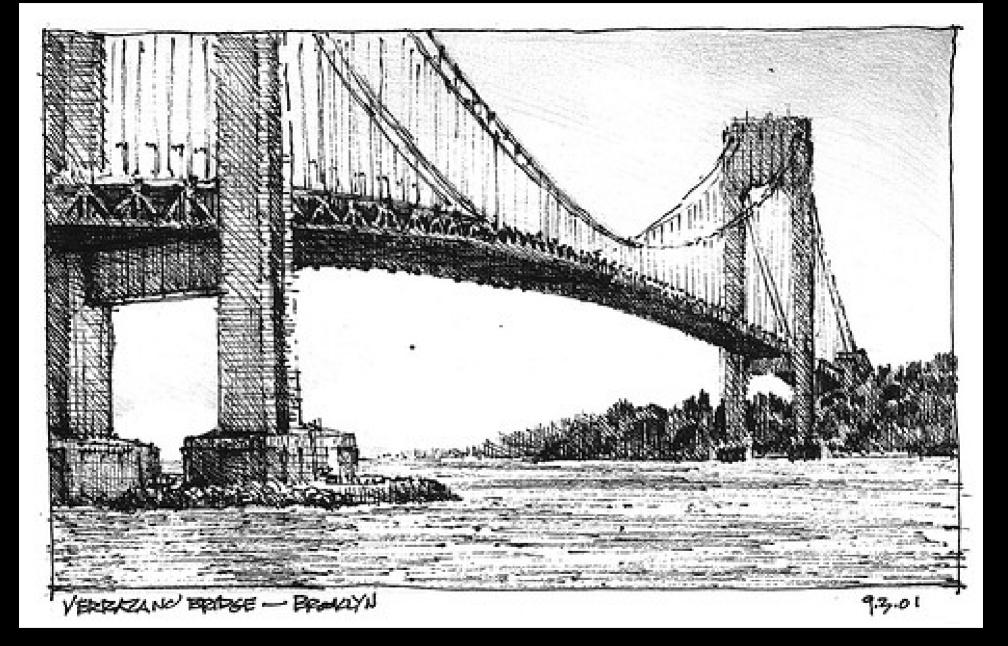


- Intercept a connection from the client side.
- Generate a certificate for the site it is connecting to.
- Sign it with any random valid leaf node certificate.
- Pass that certificate chain to the client.

- Make a normal SSL connection to the server.
- Pass data between client and server, decrypting and encrypting on each end.

LATELY, I'VE BEEN TALKING ABOUT SSL Stripping...

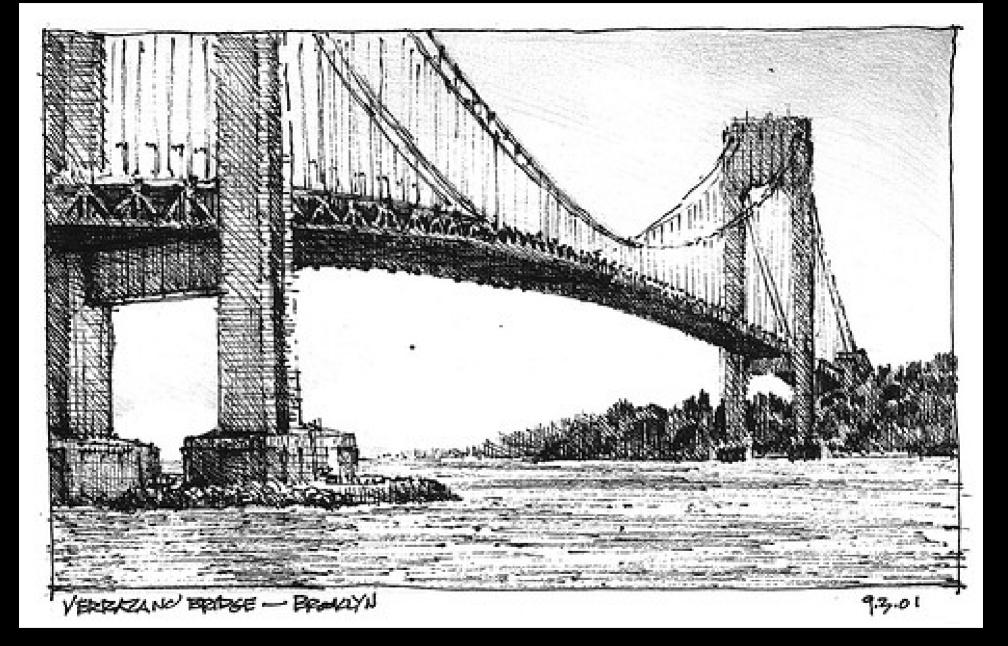
BRIEF



SSL can be useful, but how it's deployed *matters*

In the context of web browsing

- SSL is almost never encountered directly.
- It is either encountered as a result of:
 - A 302 redirect from an HTTP URL to an HTTPS URL.
 - An HTTPS link that a user clicks on from an HTTP page.
 - (Think, "My Cart," "Checkout," "Login," etc...)





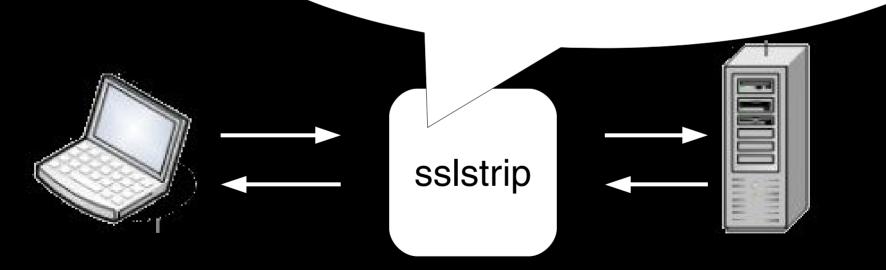
WE CAN ATTACK SSL Before We Even Get There

SSLSNIFF

SSLSTRIP

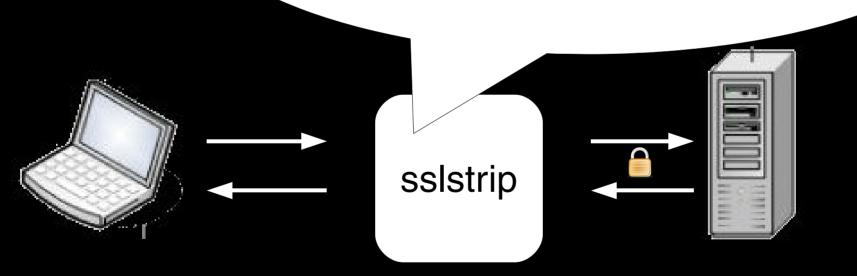
SSLSTRIP

- Watch HTTP traffic go by.
- Switch to and keep a map of what you've changed.
- Switch Location: https:// to Location: http:// and keep a map of what you've changed.



SSLSTRIP

- Watch HTTP traffic go by.
- When we see an HTTP request for a URL that we've stripped, proxy that out as HTTPS to the server.
- Watch the HTTPS traffic go by, log everything that we want, and keep a map of all relative, CSS, and JS links that go by.



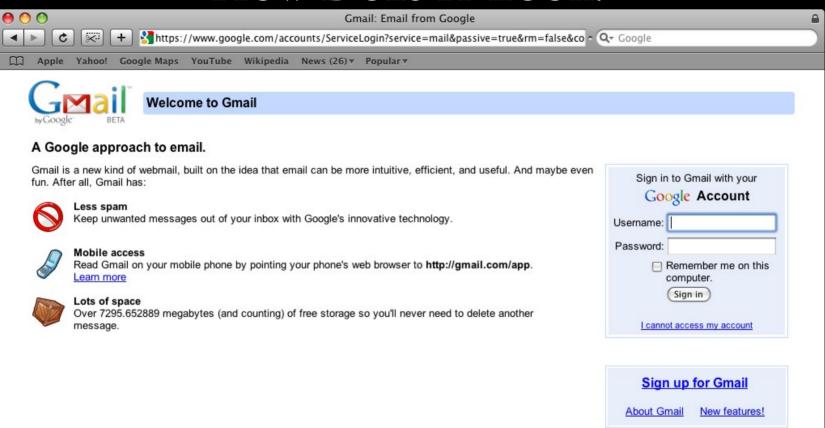
How Does It Look?

🗌 Gmail: Email from Google - Mozilla Firefox 🥄	_ = >		
<u>F</u> ile <u>E</u> dit <u>V</u> iew Hi <u>s</u> tory <u>B</u> ookmarks <u>T</u> ools <u>H</u> elp			
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<image/> Welcome to Gmail A cocyce approach to email. Brail is a new kind of webmail, built on the idea that email can be more intuitive, efficient, and useful. And maybe even fun. After all, Gmail has: Image: Stream of the strea	Sign in to Gmail with your Google Account Username: Password: Remember me on this computer. Sign in Leannot access my account		
	About Gmail New features!		
©2009 Google - Gmail for Organizations - Gmail Blog -	Terms - Help		

How Does It Look?

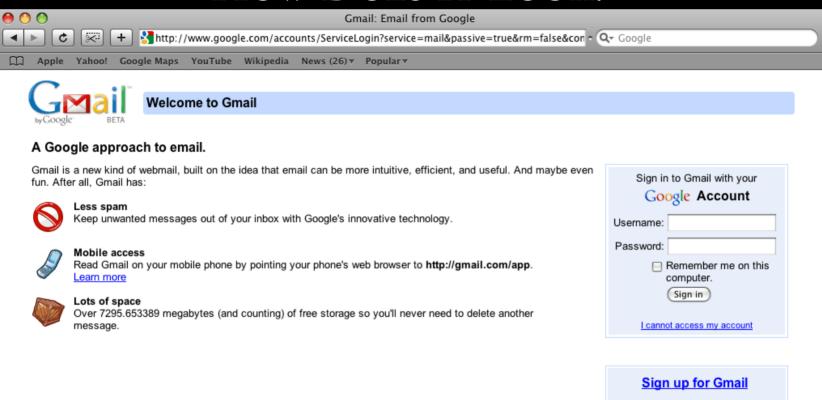
Gmail: Email from Google - Mozilla Firefox	×
<u>F</u> ile <u>E</u> dit <u>V</u> iew Hi <u>s</u> tory <u>B</u> ookmarks <u>T</u> ools <u>H</u> elp	5 ⁴ 2 *2
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🗟 Most Visited 🔻 🌘 Getting Started 🔝 Latest Headlines 🔻	
Geogle BETA Welcome to Gmail A Google approach to email.	
 Gmail is a new kind of webmail, built on the idea that email can be more intuitive, efficient, and useful. And maybe even fun. After all, Gmail has: Less spam Keep unwanted messages out of your inbox with Google's innovative technology. Mobile access Read Gmail on your mobile phone by pointing your phone's web browser to http://gmail.com/app. Learn more Lots of space Over 7290.462157 megabytes (and counting) of free storage so you'll never need to delete another message. 	Sign in to Gmail with your Google Account Username: Password: Remember me on this computer. Sign in I cannot access my account
©2009 Google - <u>Gmail for Organizations</u> - <u>Gmail Blog</u> -]	Sign up for Gmail About Gmail New features! ms - Help

How Does It Look?



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How Does It Look?



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About Gmail

New features!

Where can we go from here?

WHERE DO WE need TO GO FROM HERE?

X509Certificate

Version

Serial Number

Issuer

Validity (not before X or after Y)

Subject

PublicKey

SignatureAlgorithm

Signature

X509Certificate

Version

Serial Number

Issuer

Validity (not before X or after Y)

Subject

PublicKey

SignatureAlgorithm

Signature

X509Certificate

Version

Serial Number

Issuer

Validity (not before X or after Y)

Subject

PublicKey

SignatureAlgorithm

Signature

X509Certificate

Version

Serial Number

Issuer

Validity (not before X or after Y) Subject PublicKey SignatureAlgorithm Signature

X509Certificate

Version

Serial Number

Issuer

Validity (not before X or after Y) Subject PublicKey SignatureAlgorithm



Certificate: Data: Version: 3 (0x2) Serial Number: 01:2a:39:76:0d:3f:4f:c9:0b:e7:bd:2b:cf:95:2e:7a Signature Algorithm: shalWithRSAEncryption Issuer: C=ZA, O=Thawte Consulting (Pty) Ltd., CN=Thawte SGC CA Validity Not Before: Mar 27 22:20:07 2009 GMT Not After : Mar 27 22:20:07 2010 GMT Subject: C=US, ST=California, L=Mountain View, 0=Google Inc, CN=www.google.com Subject Public Key Info: Public Key Algorithm: rsaEncryption RSA Public Key: (1024 bit) Modulus (1024 bit): 00:d6:b9:e1:ad:b8:61:0b:1f:4e:b6:3c:09:3d:ab: e8:e3:2b:b6:e8:a4:3a:78:2f:d3:51:20:22:45:95: d8:00:91:33:9a:a7:a2:48:ea:30:57:26:97:66:c7: 5a:ef:f1:9b:0c:3f:e1:b9:7f:7b:c3:c7:cc:af:9c: d0:1f:3c:81:15:10:58:fc:06:b3:bf:bc:9c:02:b9: 51:dc:fb:a6:b9:17:42:e6:46:e7:22:cf:6c:27:10: fe:54:e6:92:6c:0c:60:76:9a:ce:f8:7f:ac:b8:5a: 08:4a:dc:b1:64:bd:a0:74:41:b2:ac:8f:86:9d:1a: de:58:09:fd:6c:0a:25:e0:79 Exponent: 65537 (0x10001) X509v3 extensions: X509v3 Extended Key Usage: TLS Web Server Authentication, TLS Web Client Authentication, Netscape Server Gated Crypto X509v3 CRL Distribution Points: URI:http://crl.thawte.com/ThawteSGCCA.crl Authority Information Access: OCSP - URI:http://ocsp.thawte.com CA Issuers - URI:http://www.thawte.com/repository/Thawte SGC CA.crt X509v3 Basic Constraints: critical CA: FALSE Signature Algorithm: shalWithRSAEncryption 39:b6:fb:11:bc:33:2c:c3:90:48:e3:6e:c3:9b:38:b1:42:d1: 00:09:58:63:a0:e1:98:1c:85:f2:ef:10:1d:60:4e:51:09:62: f5:05:bd:9d:4f:87:6c:98:72:07:80:c3:59:48:14:e2:d6:ef: d0:8f:33:6a:68:31:fa:b7:bb:85:cc:f7:c7:47:7b:67:93:3c: c3:16:51:9b:6f:87:20:fd:67:4c:2b:ea:6a:49:db:11:d1:bd: d7:95:22:43:7a:06:7b:4e:f6:37:8e:a2:b9:cf:1f:a5:d2:bd: 3b:04:97:39:b3:0f:fa:38:b5:af:55:20:88:60:93:f2:de:db:

The Big Three

Secrecy
Authenticity
Integrity



SSL/TLS Handshake Beginnings



ServerHello, ServerCertificate

SSL Handshake Beginnings

				X509Certificate
PayPai, Inc. (U	JS) https://www.paypal.com/			ADUBUEI IIIICAIE
				Version
Server Set	tings			Serial Number
	IMAP Mail Server	Det 000 Defe	lh 00	Issuer
	imap.mail.com moxie	Port: 993 Defau	ilt: 99	Validity
				Subject
				PublicKey
Login Option	ns			SignatureAlgorithm
Protocol:	IRC 🗘			Cignata or ligorithm
<u>U</u> sername	: moxiem			Signature
<u>S</u> erver:	irc.freenode.net			Ŭ

THE PROBLEMS FOR US BEGIN Attacker ClientHello ServerHello, ServerCertificate?

Let's start by looking back once more.

IN 2000, THINGS WERE DIFFERENT.

NOTARIES!

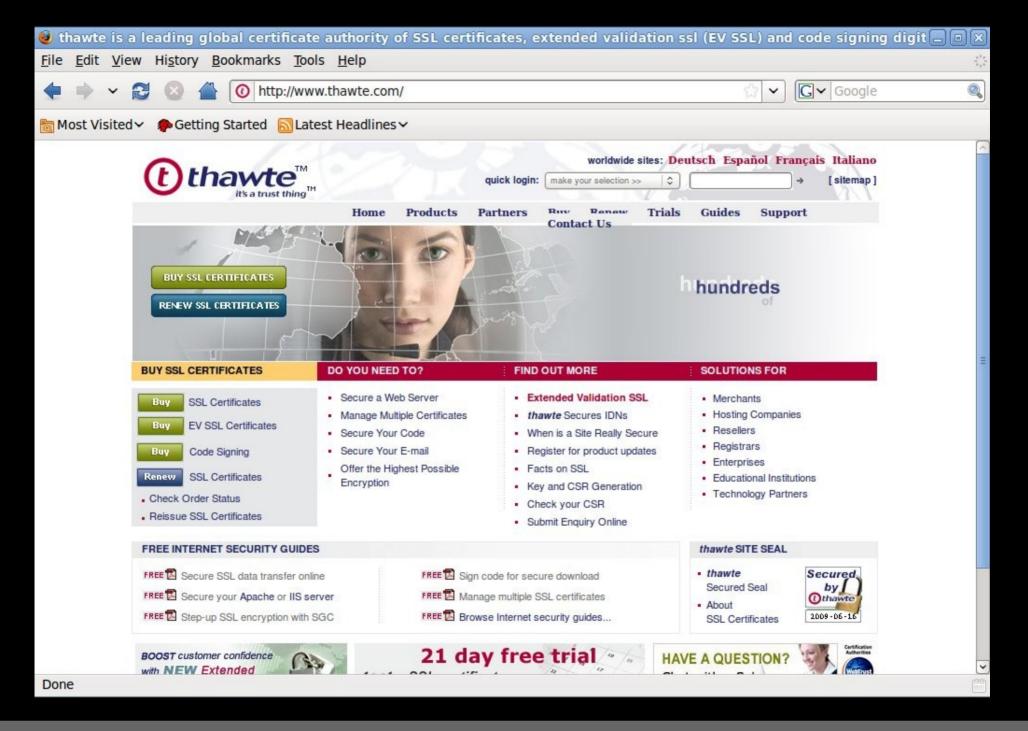
Identification!

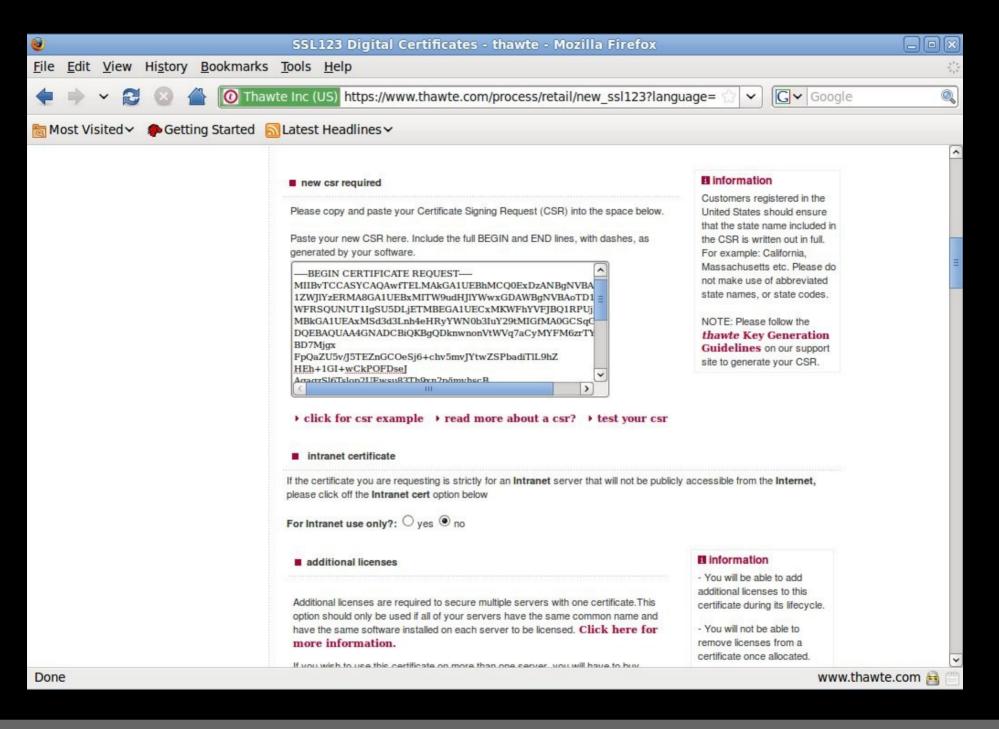
PHONE CALLS!

ACTUAL PEOPLE INVOLVED...

That is a bygone era

These days it's all about: online domain validation





CertificateRequest Version Subject PublicKey Attributes

CertificateRequest
Version
Subject
PublicKey
Attributes

CertificateRequest	
Version	
Subject	www.bankofamerica.com
PublicKey	
Attributes	

CertificateRequest	
Version	
Subject	www. <mark>bankofamerica.com</mark>
PublicKey	
Attributes	

CertificateRequest	
Version	
Subject	www. <mark>bankofamerica.com</mark>
PublicKey	
Attributes	WHOIS Lookup

CertificateRequest	
Version	
Subject	> www. <mark>bankofamerica.com</mark>
PublicKey	
Attributes	WHOIS Lookup
	Email admin@bankofamerica.com

CertificateRequest	
Version	
Subject	www <mark>.bankofamerica.com</mark>
PublicKey	
Attributes	

CertificateRequest	
Version	
Subject	
PublicKey	
Attributes	

CertificateRequest
Version
Subject
PublicKey
Attributes

certificate.authoritie s.are.a.total.ripoff. bankofamerica.com

CertificateRequest	
Version	
Subject	certificate.authoritie
PublicKey	s.are.a.total.ripoff. bankofamerica.com
Attributes	

Subjects

DistinguishedName Country State Locale Organization **Organizational Unit Common Name**

Subjects

DistinguishedName

Country State Locale Organization Organizational Unit

Common Name

- The X.509 standard is a total nightmare.
- Three revisions, twenty years.
- Parts of the standard have literally been "lost" and then later "found" again.

Subjects

DistinguishedName

Country State Locale Organization Organizational Unit Common Name

- The original vision for the DN was that each DN would fit into some global Directory Information Tree.
- In practice, the standard is weak, everyone does everything differently, and the global DIT never materialized.

Subjects

DistinguishedName

Country State Locale Organization Organizational Unit Common Name "There is nothing in any of these standards that would prevent me from including a 1 gigabit MPEG movie of me playing with my cat as one of the RDN components of the DN in my certificate."

-- Bob Jueneman on IETF-PKIX

Subjects

DistinguishedName Country State Locale Organization **Organizational Unit Common Name**

www.bankofamerica.com

CN

commonName ::= SEQUENCE { { 2 5 4 3 }, StringType(SIZE(1...64)) }

- IA5String:
 - 0x16 ID
 - 0x05 Length (5 Chars)
 - 0x76, 0x61, 0x6c, 0x75, 0x65 v, a, l, u, e

CN Encoding

• Essentially, the CN field is represented as a PASCAL String.

0xe	W	W	W	•	р	а	у	р	а	I		с	0	m
-----	---	---	---	---	---	---	---	---	---	---	--	---	---	---

• This is different from how C strings are represented.

W	w	W		р	а	у	р	а	I	•	С	0	m	\0
---	---	---	--	---	---	---	---	---	---	---	---	---	---	----

DistinguishedNa	ame		
Country			
State			
Locale			
Organization	7		
Organizational	t /nit ght	crime.org	
Common Name		— WWW.	paypal.com

Common Name

www.thoughtcrime.org

Common Name

verisign.eats.children.thoughtcrime.org

Common Name

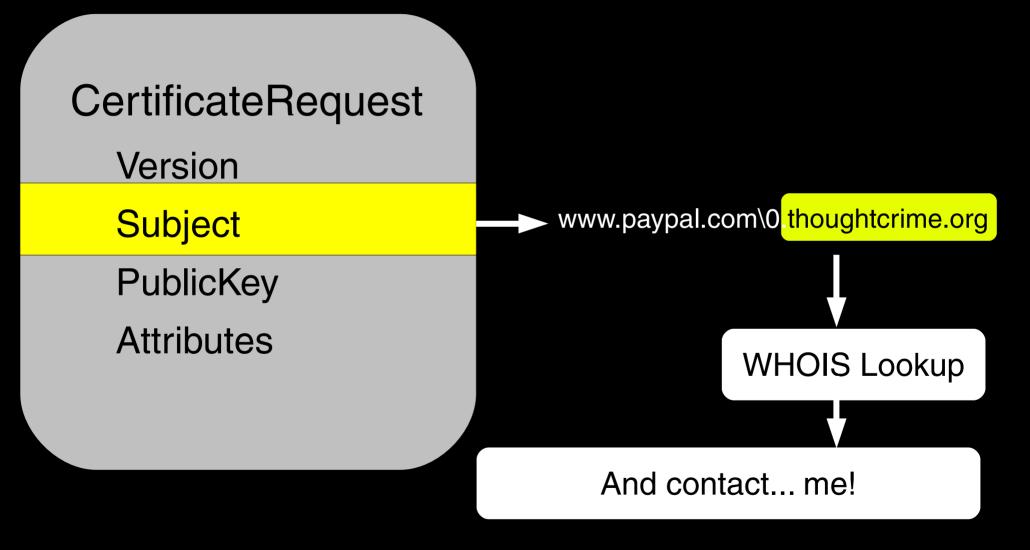
Common Name

www.paypal.com\0.thoughtcrime.org

PKCS #10 Certificate Signing Request

CertificateRequest	
Version	
Subject	www.paypal.com\0.thoughtcrime.org
PublicKey	
Attributes	

PKCS #10 Certificate Signing Request



Our Original Scenario





Our Original Scenario

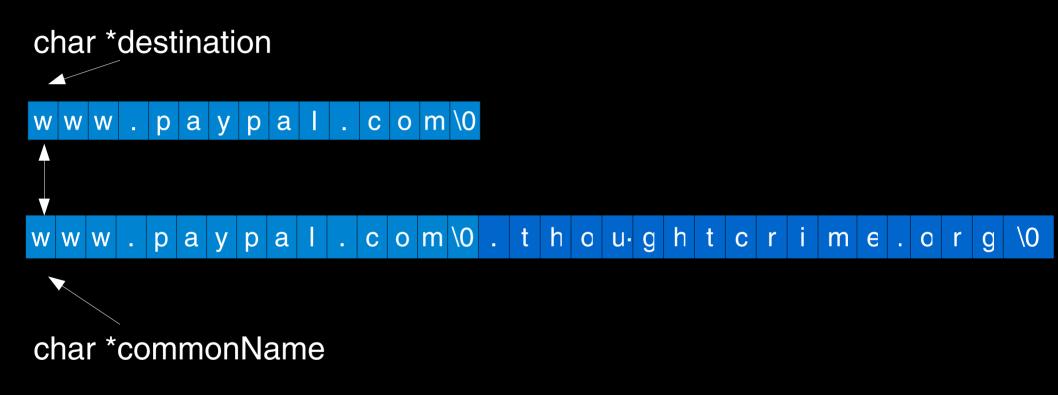


X509Certificate Version Serial Number Issuer Validity Subject PublicKey SignatureAlgorithm Signature

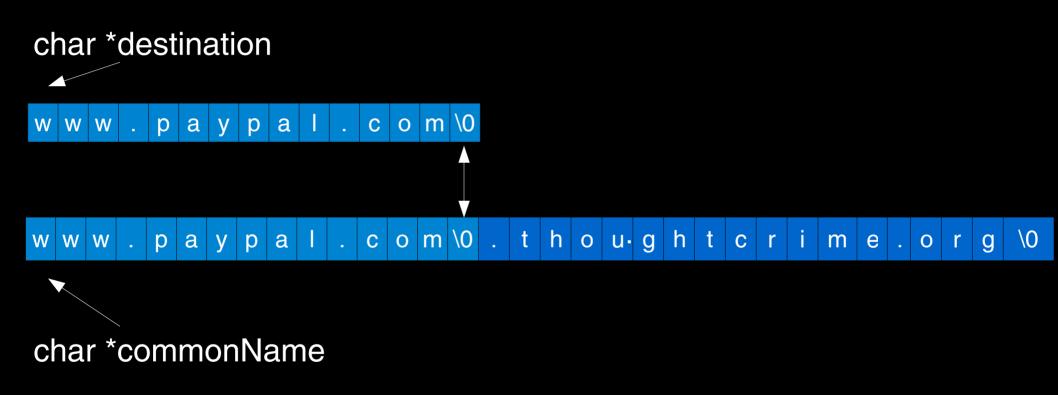
Our Original Scenario

char *destination = getDomainWeAreConnectingTo(); char *commonName = getCommonNameFromCertificate(); bool everythingIsOk = (strcmp(destination, commonName) == 0);

IN MEMORY, THOUGH...



IN MEMORY, THOUGH...

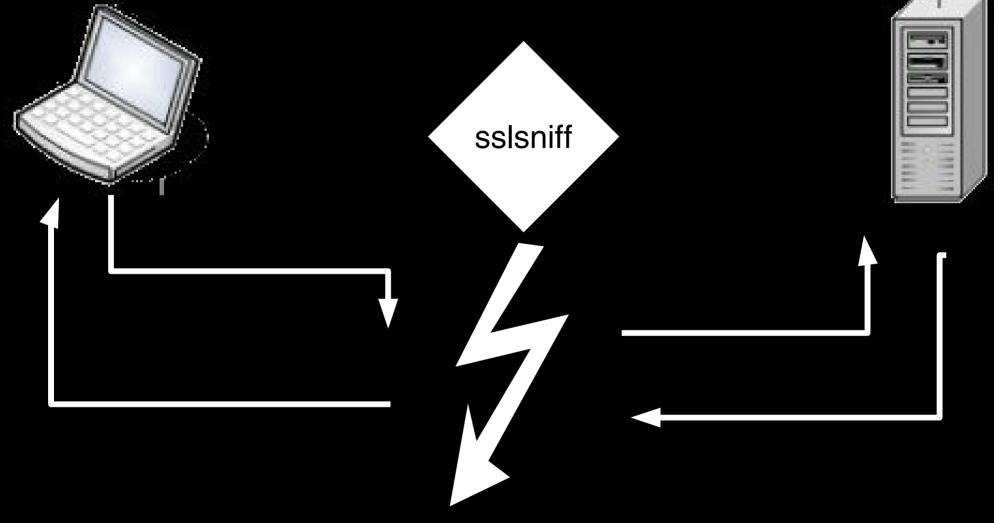


IN THE EYES OF MOST SSL IMPLEMENTATIONS, THIS CERTIFICATE IS *completely valid* FOR WWW.PAYPAL.COM

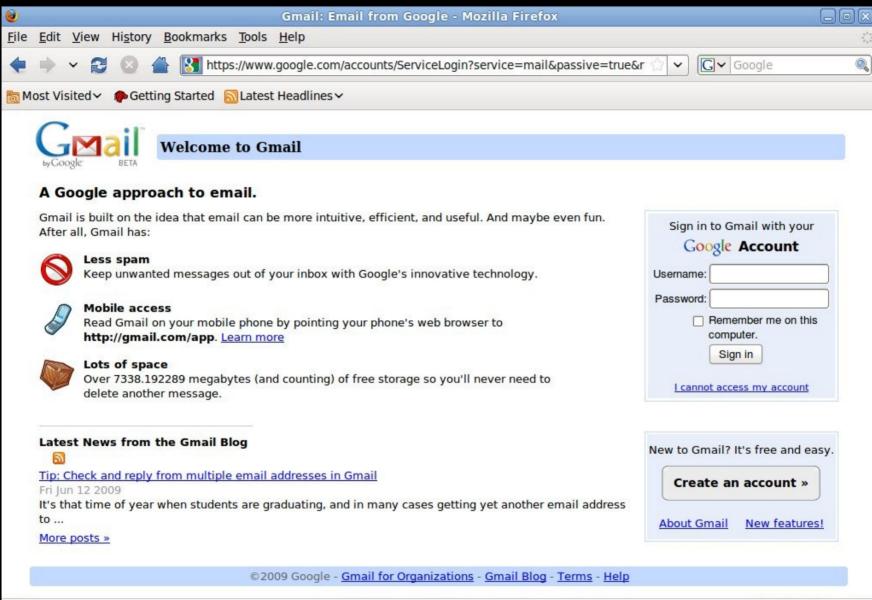
What are "most" SSL implementations?

- Web Browsers
 - Firefox (all versions), IE (all versions), Lynx, Curl,
- Mail Clients
 - Thunderbird, Outlook, Evolution
- Chat Clients
 - Pidgin, AIM, irssi, centericq
- SSL VPNs
 - AEP, Citrix, etc...

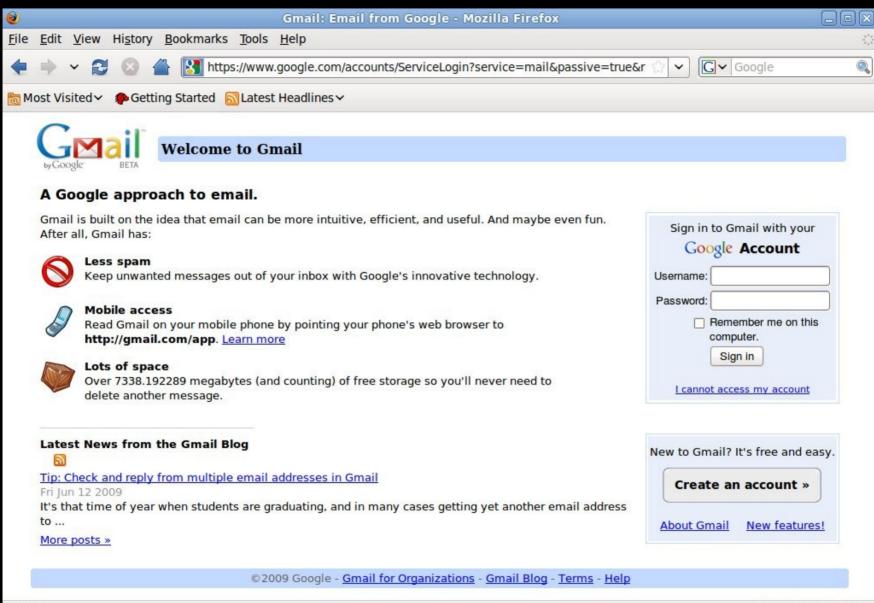
A First Cut: updated sslsniff



Iff "null prefix attack" certificate is available



www.google.com 🔒



www.google.com

Moxie Marlinspike Institute For Disruptive Studies

Done

General Media Permissions Security		
Web Site Identity		
Web site: www.google.com		
Owner: This web site does not su	upply identity i	nformation.
Verified by:	-	
-		
This web site provides a certificate to verify	its identity.	View Certificate
Privacy & History		
Have I visited this web site before today?	Yes, 9 times	6 - C C C C C C C C
Is this web site storing information (cookies) on my computer?	Yes	View Cookies
Have I saved any passwords for this web site?	No	Vie <u>w</u> Saved Passwords
Technical Details		
Connection Encrypted: High-grade En The page you are viewing was encrypted be		
Encryption makes it very difficult for unauth traveling between computers. It is therefore as it traveled across the network.		

Page Info - https://www.google.com/ac	counts/Service	Login?service 🗕 🗖
General Media Permissions Security		
Web Site Identity		
Web site: www.google.com		
Owner: This web site does not su	nnlv identitv i	nformation.
Verified by:	ppij lacinity i	
venned by.		
This web site provides a certificate to verify	its identity.	View Certificate
Paire and C. Minta		
Privacy & History		
Have I visited this web site before today?	Yes, 9 times	
Is this web site storing information (cookies) on my computer?	Yes	View Cookies
Have I saved any passwords for this web site?	No	Vie <u>w</u> Saved Passwords
Technical Details		
Connection Encrypted: High-grade End	ryption (RC4 1	28 bit)
The page you are viewing was encrypted be		
Encryption makes it very difficult for unauth traveling between computers. It is therefore as it traveled across the network.	orized people to	view information

DISADVANTAGES

I) TARGETED ATTACKS ARE KIND OF LAME.

Maybe there's another trick in here somewhere...

X

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```
static int
     _shexp_match(const char *str, const char *exp, PRBool case_insensitive)
         register int x,y;
         int ret, neg;
         ret = 0;
         for(x=0,y=0;exp[y];++y,++x)
             if((!str[x]) && (exp[y] != '(') && (exp[y] != '$') && (exp[y] != '*'))
                 ret = ABORTED;
                 switch(exp[y]) {
                   case '$':
                     if( (str[x]) )
                         ret = NOMATCH;
                          --×;
                     break;
                   case '*':
                     while(exp[++y] == '*'){}
                     if(!exp[y])
                         return MATCH;
                     while(str[x]) {
                          switch(_shexp_match(&str[x++],&exp[y], case_insensitive)) {
                          case NOMATCH:
                              continue;
                         case ABORTED:
                              ret = ABORTED;
                              break;
                          default:
                              return MATCH;
                         break;
                     if((exp[y] == '$') && (exp[y+1] == '\0') && (!str[x]))
                         return MATCH;
                         ret = ABORTED;
                   break;
case '[':
                     neg = ((exp[++y] == '^') && (exp[y+1] != ']'));
                     if (neg)
-u:** *scratch*
                       6% L46
                                  (C/1 Abbrev)--
```

File Edit Options Burners Tools C meg

```
a constal in 101 as constal in 101 as constal in 1013
to Passadular
                   10.04
```

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```
static int
     _shexp_match(const char *str, const char *exp, PRBool case_insensitive)
         register int x,y;
         int ret, neg;
         ret = 0;
         for(x=0,y=0;exp[y];++y,++x) {
             if((!str[x]) && (exp[y] != '(') && (exp[y] != '$') && (exp[y] != '*'))
                 ret = ABORTED;
                 switch(exp[y]) {
                   case '$':
                     if( (str[x]) )
                         ret = NOMATCH;
                          --×;
                     break;
                   case '*':
                     while(exp[++y] == '*'){}
                     if(!exp[y])
                         return MATCH;
                     while(str[x]) {
                          switch(_shexp_match(&str[x++],&exp[y], case_insensitive)) {
                         case NOMATCH:
                              continue;
                         case ABORTED:
                              ret = ABORTED;
                              break;
                         default:
                             return MATCH;
                         break;
                     if((exp[y] == '$') && (exp[y+1] == '\0') && (!str[x]))
                         return MATCH;
                         ret = ABORTED;
                     break;
                     neg = ((exp[++y] == '^') && (exp[y+1] != ']'));
                     if (neg)
-u:** *scratch*
                       6% L46
                                  (C/1 Abbrev)-----
```

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```
static int
     _shexp_match(const char *str, const char *exp, PRBool case_insensitive)
         register int x,y;
         int ret, neg;
         ret = 0;
         for(x=0,y=0;exp[y];++y,++x) {
             if(([str[x]) && (exp[y] != '(') && (exp[y] != '$') && (exp[y] != '*'))
                 ret = ABORTED;
                 switch(exp[y]) {
                   case '$':
                     if( (str[x]) )
                         ret = NOMATCH;
                          --×;
                     break;
                   case '*':
                     while(exp[++y] == '*'){}
                     if(!exp[y])
                         return MATCH;
                     while(str[x]) {
                          switch(_shexp_match(&str[x++],&exp[y], case_insensitive)) {
                         case NOMATCH:
                              continue;
                         case ABORTED:
                              ret = ABORTED;
                              break;
                         default:
                             return MATCH;
                         break;
                     if((exp[y] == '$') && (exp[y+1] == '\0') && (!str[x]))
                         return MATCH;
                         ret = ABORTED;
                     break;
                     neg = ((exp[++y] == '^') && (exp[y+1] != ']'));
                     if (neg)
-u:** *scratch*
                       6% L46
                                  (C/1 Abbrev)-----
```

🗅 🗁 📃 🗶 🔚 🖪 🥱 🔏 🖬 🛅 🖳 🎘 🤯

```
static int
     _shexp_match(const char *str, const char *exp, PRBool case_insensitive)
         register int x,y;
         int ret, neg;
         ret = 0;
         for(x=0,y=0;exp[y];++y,++x) {
             if((!str[x]) && (exp[y] != '(') && (exp[y] != '$') && (exp[y] != '*'))
                 ret = ABORTED;
                 switch(exp[y]) {
                   case '$':
                     if( (str[x]) )
                         ret = NOMATCH;
                          --×;
                     break;
                   case '*':
                     while(exp[++y] == '*'){}
                     if(!exp[y])
                         return MATCH;
                     while(str[x]) {
                          switch(_shexp_match(&str[x++],&exp[y], case_insensitive)) {
                         case NOMATCH:
                              continue;
                         case ABORTED:
                              ret = ABORTED;
                              break;
                         default:
                             return MATCH;
                         break;
                     if((exp[y] == '$') && (exp[y+1] == '\0') && (!str[x]))
                         return MATCH;
                         ret = ABORTED;
                     break;
                     neg = ((exp[++y] == '^') && (exp[y+1] != ']'));
                     if (neg)
-u:** *scratch*
                       6% L46
                                  (C/1 Abbrev)-----
```

Universal Wildcard

*\0.thoughtcrime.org

Universal Wildcard

*~.thoughtcrime.org

Other Weird Stuff

 (www.paypal.com|mail.google.com| www.etrade.com|www.bankofamerica.com| www.wachovia.com|www.pnc.com| www.wellsfargo.com)\0.thoughtcrime.org

And... your remote exploit.

- 144 char *e2 = (char *) PORT_Alloc(sizeof(char)*strlen(exp));
- 145 register int t,p2,p1 = 1;
- 146 int cp;
- 147
- 148 while(1) {
- 149 for(cp=1;exp[cp] != ')';cp++)
- 150 if(exp[cp] == '\\')
- 151 ++cp;
- 152 for(p2 = 0;(exp[p1] != '|') && (p1 != cp);p1++,p2++) {

```
• 153 if(exp[p1] == '\\\')
```

• 154 e2[p2++] = exp[p1++];

}

• 156

```
• 157 for (t=cp+1; ((e2[p2] = exp[t]) != 0); ++t,++p2)
```

• 158 if(_shexp_match(str,e2, case_insensitive) == MATCH) {

```
• 159 PORT_Free(e2);
```

```
• 160 return MATCH;
```

- 161 }
- 162

And... your remote exploit.

- . 144 char *e2 = (char *) PORT_Alloc(sizeof(char)*strlen(exp));
- 145 register int t,p2,p1 = 1;
- 146 int cp;
- 147
- 148 while(1) {
- 149 for(cp=1;exp[cp] != ')';cp++)
- 150 if(exp[cp] == '\\')
- 151 ++cp;

}

- 152 for(p2 = 0;(exp[p1] != '|') && (p1 != cp);p1++,p2++) {
- 153 if(exp[p1] == '\\')
- 154 e2[p2++] = exp[p1++];
- 155 e2[p2] = exp[p1];
- 156
- 157 for (t=cp+1; ((e2[p2] = exp[t]) != 0); ++t,++p2) {}
- 158 if(_shexp_match(str,e2, case_insensitive) == MATCH) {

```
• 159 PORT_Free(e2);
```

- 160 return MATCH;
- 161 }
- 162

- . 144 char *e2 = (char *) PORT_Alloc(sizeof(char)*strlen(exp));
- 145 register int t,p2,p1 = 1;
- 146 int cp;
- 147
- 148 while(1) {
- 149 for(cp=1;exp[cp] != ')';cp++)
- 150 if(exp[cp] == '\\')
- 151 ++cp;

}

- 152 for(p2 = 0;(exp[p1] != '|') && (p1 != cp);p1++,p2++) {
- 153 if(exp[p1] == '\\')
- 154 e2[p2++] = exp[p1++];
- 155 e2[p2] = exp[p1];
- 156
- 157 for (t=cp+1; ((e2[p2] = exp[t]) != 0); ++t,++p2)
- 158 if(_shexp_match(str,e2, case_insensitive) == MATCH) {

```
• 159 PORT_Free(e2);
```

- 160 return MATCH;
- 161 }
- 162 ...

- . 144 char *e2 = (char *) PORT_Alloc(sizeof(char)*strlen(exp));
- 145 register int t,p2,p1 = 1;
- 146 int cp;
- 147
- 148 while(1) {
- 149 for(cp=1;exp[cp] != ')';cp++)
- 150 if(exp[cp] == '\\')
- 151 ++cp;

}

- 152 for(p2 = 0;(exp[p1] != '|') && (p1 != cp);p1++,p2++) {
- 153 if(exp[p1] == '\\')
- 154 e2[p2++] = exp[p1++];
- 155 e2[p2] = exp[p1];
- 156
- 157 for (t=cp+1; ((e2[p2] = exp[t]) != 0); ++t,++p2) {}
- 158 if(_shexp_match(str,e2, case_insensitive) == MATCH) {
- 159 PORT_Free(e2);
- 160 return MATCH;
- 161 }
- 162

AND... YOUR REMOTE EXPLOIT.

No signed signature required!

Possible to sneak non-ASCII characters past the NSS filters.

This yields something exploitable in Firefox, Thunderbird, Evolution, Pidgin, and AIM.

A SECOND CUT: SSLSNIFF WITH WILDCARD SUPPORT sslsniff

- Perform MITM if "null termination attack" cert is available.
- Or perform MITM with "universal wildcard" cert if client is NSS.

A Second Cut: updated sslsniff

sslsniff

- Watches network and fingerprints clients for level of vulnerability.
- Every NSS client's communication is intercepted – either with a specific "null termination" certificate, or with the "universal wildcard" certificate.
- Every non-NSS client that is vulnerable is intercepted with a "null termination" certificate if available for the destination host.
- Non-vulnerable clients are left alone to avoid detection.

WHAT DO WE HAVE TO WORRY ABOUT?

1) Certificate Revocation



1) Certificate Revocation

 It would be unfortunate if some bitter Certificate Authority decided to revoke our universal wildcard certificates or any of our nulltermination certificates.

2) Updates

 It would be unfortunate if some bitter SSL implementation decided to start paying attention to how ASN.1 is formatted.

1) Certificate Revocation

- These days, it's all about Online Certificate Status Protocol (OCSP).
- Whenever a SSL stack sees a new certificate, it makes a quick request to the OCSP URL that the signing CA embedded in it.
- The SSL stack receives a signed response from the OCSP provider indicating whether the certificate has been revoked or not.

Certificate: Data: Version: 3 (0x2) Serial Number: 01:2a:39:76:0d:3f:4f:c9:0b:e7:bd:2b:cf:95:2e:7a Signature Algorithm: shalWithRSAEncryption Issuer: C=ZA, O=Thawte Consulting (Pty) Ltd., CN=Thawte SGC CA Validity Not Before: Mar 27 22:20:07 2009 GMT Not After : Mar 27 22:20:07 2010 GMT Subject: C=US, ST=California, L=Mountain View, O=Google Inc, CN=www.google.com Subject Public Key Info: Public Key Algorithm: rsaEncryption RSA Public Key: (1024 bit) Modulus (1024 bit): 00:d6:b9:e1:ad:b8:61:0b:1f:4e:b6:3c:09:3d:ab: e8:e3:2b:b6:e8:a4:3a:78:2f:d3:51:20:22:45:95: d8:00:91:33:9a:a7:a2:48:ea:30:57:26:97:66:c7: 5a:ef:f1:9b:0c:3f:e1:b9:7f:7b:c3:c7:cc:af:9c: d0:1f:3c:81:15:10:58:fc:06:b3:bf:bc:9c:02:b9: 51:dc:fb:a6:b9:17:42:e6:46:e7:22:cf:6c:27:10: fe:54:e6:92:6c:0c:60:76:9a:ce:f8:7f:ac:b8:5a: 08:4a:dc:b1:64:bd:a0:74:41:b2:ac:8f:86:9d:1a: de:58:09:fd:6c:0a:25:e0:79 Exponent: 65537 (0x10001) X509v3 extensions: X509v3 Extended Key Usage: TLS Web Server Authentication, TLS Web Client Authentication, Netscape Server Gated Crypto X509v3 CRL Distribution Points: URI:http://crl.thawte.com/ThawteSGCCA.crl Authority Information Access: OCSP - URI:http://ocsp.thawte.com CA Issuers - URI:http://www.thawte.com/repository/Thawte SGC CA.crt X509v3 Basic Constraints: critical CA: FALSE Signature Algorithm: shalWithRSAEncryption 39:b6:fb:11:bc:33:2c:c3:90:48:e3:6e:c3:9b:38:b1:42:d1: 00:09:58:63:a0:e1:98:1c:85:f2:ef:10:1d:60:4e:51:09:62: f5:05:bd:9d:4f:87:6c:98:72:07:80:c3:59:48:14:e2:d6:ef: d0:8f:33:6a:68:31:fa:b7:bb:85:cc:f7:c7:47:7b:67:93:3c: c3:16:51:9b:6f:87:20:fd:67:4c:2b:ea:6a:49:db:11:d1:bd: d7:95:22:43:7a:06:7b:4e:f6:37:8e:a2:b9:cf:1f:a5:d2:bd: 3b:04:97:39:b3:0f:fa:38:b5:af:55:20:88:60:93:f2:de:db:

OCSPResponse ::= SEQUENCE {

- responseStatus OCSPResponseStatus,
- responseBytes

}

[0] EXPLICIT ResponseBytes OPTIONAL

```
OCSPResponse ::= SEQUENCE {
                     OCSPResponseStatus,
   responseStatus
                     [0] EXPLICIT ResponseBytes OPTIONAL
   responseBytes
ResponseBytes ::=
                   SEQUENCE {
   responseType OBJECT IDENTIFIER,
              OCTET STRING
   response
BasicOCSPResponse ::= SEQUENCE {
   tbsResponseData
                     ResponseData,
   signatureAlgorithm AlgorithmIdentifier,
   signature
                 BIT STRING,
                [0] EXPLICIT SEQUENCE OF Certificate OPTIONAL }
   certs
```

```
OCSPResponse ::= SEQUENCE {
                     OCSPResponseStatus,
   responseStatus
                     [0] EXPLICIT ResponseBytes OPTIONAL
   responseBytes
ResponseBytes ::=
                   SEQUENCE {
   responseType OBJECT IDENTIFIER,
               OCTET STRING
   response
BasicOCSPResponse ::= SEQUENCE {
   tbsResponseData
                      ResponseData,
   signatureAlgorithm AlgorithmIdentifier,
   signature
                 BIT STRING.
                [0] EXPLICIT SEQUENCE OF Certificate OPTIONAL }
   certs
```

OCSPResponse ::= SEQUENCE {		
responseSta	tus OCSPRespo	onseStatus,
responseByte	es [0] EXPLICI	ResponseBytes OPTIONAL
}		
ResponseBytes :	:= SEQUENCE {	
responseTyp	e OBJECT IDENTIF	FIER,
response	OCTET STRING	
}		
BasicOCSPResp	oonse ::= SEQUE	
tbsResponse	Data ResponseD	ata,
signatureAlg	orithm AlgorithmIde	ntifier,
signature	BIT STRING,	
certs	[0] EXPLICIT SEQ	UENCE OF Certificate OPTIONAL }

OCSPResponse ::= SEQUENCE {

responseStatus	OCSPResponseStatus,
responseBytes	[0] EXPLICIT ResponseBytes OPTIONAL
}	
ResponseBytes ::=	SEQUENCE {
responseType (OBJECT IDENTIFIER,
response OC	CTET STRING
}	
BasicOCSPRespons	se ::= SEQUENCE {
tbsResponseDat	ta ResponseData,
signatureAlgorith	nm AlgorithmIdentifier,
signature	BIT STRING,
certs [0)] EXPLICIT SEQUENCE OF Certificate OPTIONAL }

OCSPResponse ::= SEQUENCE {	
responseStatus OCSPResponseStatus,	
responseBytes [0] EXPLICIT ResponseBytes OPTIONAL	
}	
ResponseBytes ::= SEQUENCE {	
responseType OBJECT IDENTIFIER,	
response OCTET STRING	
}	
BasicOCSPResponse ::= SEQUENCE {	
tbsResponseData ResponseData,	
signatureAlgorithm AlgorithmIdentifier,	
signature BIT STRING,	
certs [0] EXPLICIT SEQUENCE OF Certificate OPTIONAL }	

OCSPResponse ::=	= SEQUENCE {
responseStatus	s OCSPResponseStatus,
responseBytes	[0] EXPLICIT ResponseBytes OPTIONAL
}	
OCSPResponseSt	atus ::= ENUMERATED {
successful	(0),Response has valid confirmations
malformedReq	uest (1),Illegal confirmation request
internalError	(2),Internal error in issuer
tryLater	(3),Try again later
	(4) is not used
sigRequired	(5),Must sign the request
unauthorized	(6)Request unauthorized

OCSPResponse ::= SEQUENCE {			
responseStatus	oCSPResponseStatus,		
responseBytes	[0] EXPLICIT ResponseBytes OPTIONAL		
}			
OCSPResponseStatus ::= ENUMERATED {			
successful	(0),Response has valid confirmations		
malformedRequ	uest (1),Illegal confirmation request		
internalError	(2),Internal error in issuer		
tryLater	(3),Try again later		
	(4) is not used		
sigRequired	(5),Must sign the request		
unauthorized	(6)Request unauthorized		
}			

OCSPResponse ::= SEQUENCE {

responseStatus OCSPResponseStatus = 3,

responseBytes

}

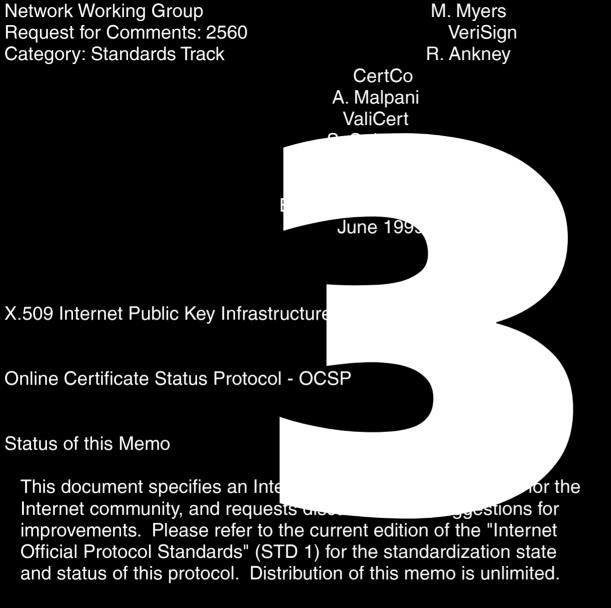
[0] EXPLICIT ResponseBytes OPTIONAL

OCSPResponse ::= SEQUENCE {

}

responseStatus OCSPResponseStatus = 3,

PROPOSED STANDARD



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A Third Cut: ocsp-aware sslsniff

sslsniff

- Watch network and fingerprints clients for level of vulnerability.
- Every NSS client's communication is intercepted either with a specific "null termination" certificate, or with the "universal wildcard" certificate.
- Every non-NSS client that is vulnerable is intercepted with a "null termination" certificate if available for the destination host.
- Non-vulnerable clients are left alone to avoid detection.
- Optionally watch for OCSP requests corresponding to certificates we're using, and "tryLater" them to defeat OCSP.

2) Updates

- It used to be that people, you know, downloaded and installed updates.
- As software gets more complicated, it is inevitably shipped with more bugs, and attackers are situated to exploit them on a larger scale.
- So some have felt the need to deploy self-updating software in order to fix problems rapidly.

2) Updates

- This is bad news for us, because by standing here and talking to you about this stuff, it probably means that SSL implementations are going to fix these problems.
- But their update mechanisms in themselves seem like kind of a dangerous idea, right?
- Maybe there's something we can do about our problem.

- When you install Firefox, it comes with a feature called "automatic update service," which happens to be enabled by default.
- Here be dragons.

Update Server In The Sky

Connection to

S

aus2.mozilla.org



Hello, do you have any updates for me? Here's my product, version, build ID, OS, locale, and channel. As a matter of fact, I do. Here's an unsigned blob of data – you'd do well to install it.

- Firefox and Thunderbird depend on their TLS connection to the update server to defend them against all possible attacks.
- Code is returned from the update server either as a binary diff against the distribution binary the client is running, or as a complete image of the binary.
- By default, "minor" updates are downloaded and installed silently – only prompting the user to restart their browser once everything is done.
 - The update server is the one who reports the version number of the update, so it is effectively up to the server whether the image it provides is installed silently or not.

- As vendors start to release patches for this vulnerability, the update mechanisms themselves will be vulnerable.
- All we need is a universal wildcard cert, or alternately a null-termination prefix cert for aus2.mozilla.org, and we can take control of the update mechanism to deliver payloads of our choice.
 - This could be anything:
 - A rootkit that logs keystrokes.
 - Something that sends all traffic/email through a server of our choosing.
 - A completely legitimate image that just happens to include our own CA certs.
 - Or, just to be confusing, a totally different web browser ("Thank you for updating to Galeon 0.0.3!") or even a completely different type of application – notepad.exe comes to mind.

 In order to patch your system effectively, you will not be able to trust anything that comes through automatic updates.

A Fourth Cut: update-aware sslsniff

 Watch network and fingerprints clients for level of vulnerability.

sslsniff

- Every NSS client's communication is intercepted either with a specific "null prefix" certificate, or with the "universal wildcard" certificate.
- Every non-NSS client that is vulnerable is intercepted with a "null prefix" certificate if available for the destination host.
- Non-vulnerable clients are left alone to avoid detection.
- Optionally watch for OCSP requests corresponding to certificates we're using, and "tryLater" them to defeat OCSP.
- Optionally watch for Firefox/Thunderbird update polls, and respond with a "custom" build.

Postscript: Stripping NULL is no solution

- Some SSL/TLS implementations (Safari, Opera) appear to strip '\0' from commonName strings before comparing.
- Thus:

www.paypal.com\0.thoughtcrime.org

Becomes:

www.paypal.com.thoughtcrime.org

Postscript: Stripping NULL is no solution

- These implementations are vulnerable to a variation of our attack.
- The key is that some Certificate Authorities are vulnerable to this attack internally.
 - When presented with www.paypal.com\0.thoughtcrime.org, some CAs internally validate it as www.paypal.com
 - But the whole string (www.paypal.com\0.thoughtcrime.org) is what ends up in the subject of the cert they later issue.

Postscript: Stripping NULL is no solution

- So if we register a domain like sitekey.ba
- We can get a certificate for sitekey.ba\Onkofamerica.com
- The CAs that are internally vulnerable to this attack will validate that certificate against sitekey.ba, which we own.
- When the cert is later presented to a SSL implementation that strips \0, the certificate's common name becomes: sitekey.bankofamerica.com

Conclusion

- We have a MITM attack that will intercept communication for almost all SSL/TLS implementations.
- In the case of NSS (Firefox, Thunderbird, Evolution, AIM, Pidgin) we only need a single certificate.
- We've defeated the OCSP protocol as implemented.
- We've hijacked the Mozilla auto-updates for both applications and extensions.
- We've got an exploitable overflow.
- In short, we've got your passwords, your communication, and control over your computer.